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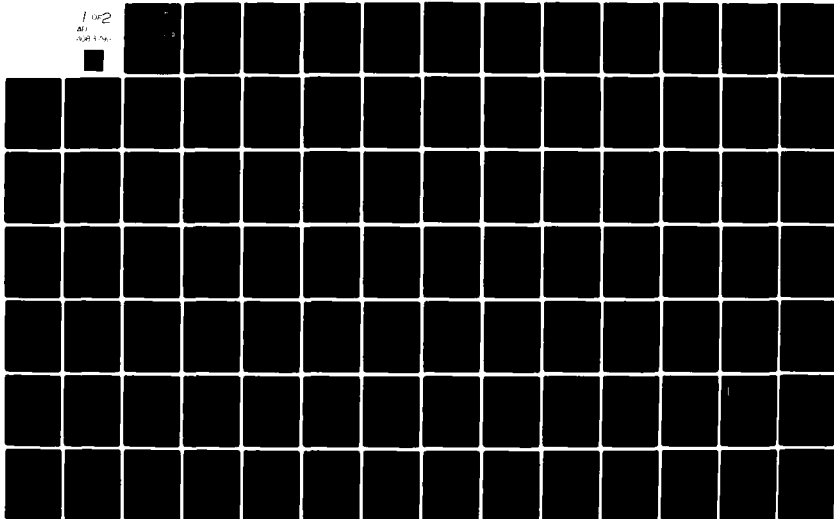
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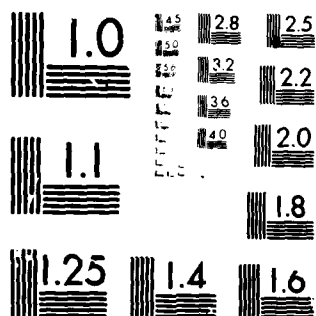
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An Analysis of the Multiple Objective Capital
Budgeting Problem Via Fuzzy Linear Integer
(0-1) Programming

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May 31, 1980

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A thesis submitted to The Pennsylvania State University,
University Park, Pennsylvania, in partial fulfillment
of the requirements for the degree of Master of Science
in Industrial Engineering and Operations Research.

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REPORT DOCUMENTATION PAGE		READ INSTRUCTIONS BEFORE COMPLETING FORM
1. REPORT NUMBER	2. GOVT ACCESSION NO.	3. RECIPIENT'S CATALOG NUMBER
	AD-A083 792	
4. TITLE (and Subtitle)		5. TYPE OF REPORT & PERIOD COVERED
6 An Analysis of the Multiple Objective Capital Budgeting Problem via Fuzzy Linear Integer (0-1) Programming.		7 Final Report, 31 May 1980
7. AUTHOR(s)		8. PERFORMING ORG. REPORT NUMBER
CPT. Michael G. Headly 10		
9. PERFORMING ORGANIZATION NAME AND ADDRESS		10. CONTRACT OR GRANT NUMBER(s)
Student, HQDA, MILPERCEN, ATTN: DAPC-CPP-E 200 Stovall Street, Alexandria, Va. 22332		
11. CONTROLLING OFFICE NAME AND ADDRESS		12. REPORT DATE
HQDA, MILPERCEN, ATTN: DAPC-OPP-E 200 Stovall Street, Alexandria VA 22332		11 31 May 1980
14. MONITORING AGENCY NAME & ADDRESS (if different from Controlling Office)		13. NUMBER OF PAGES
12 154		
15. SECURITY CLASS. (of this report)		15a. DECLASSIFICATION/DOWNGRADING SCHEDULE
Unclassified		
16. DISTRIBUTION STATEMENT (of this Report)		
Approved for Public Release; Distribution Unlimited.		
17. DISTRIBUTION STATEMENT (of the abstract entered in Block 20, if different from Report)		
18. SUPPLEMENTARY NOTES		
Master's of Science Thesis Submitted to the Pennsylvania State University's Industrial Engineering and Operations Research Program. Portions of this thesis were presented at the TMS/CRSA Joint National Meeting, Washington D.C., May 1980.		
19. KEY WORDS (Continue on reverse side if necessary and identify by block number)		
Linear Programming, Integer Programming, Capital Budgeting, Project and Program Selection, Heuristic Solution, Research and Development, Optimization, Mathematical Programming.		
20. ABSTRACT (Continue on reverse side if necessary and identify by block number)		
A multiple objective fuzzy linear programming approach to the capital budgeting problem is developed. Since much of the available data in any capital budgeting decision situation is either of an imprecise or ill-defined nature, a mathematical optimization technique is required that is capable of incorporating this inherent uncertainty. Fuzzy linear programming provides an effective methodology for this analysis. Specifically, a mathematical model is developed which utilizes		

20. Continued

fuzzy linear programming as a solution technique for the research and development program or project selection problem. In addition an exchange heuristic, a modified form of C. C. Petersen's exchange algorithm, is presented.

A limited bibliography of works in multiple objective optimization is presented. Two computer codes are included. The first utilizes the IBM MPSX/Mixed Integer Programming procedures to solve the (0-1) linear integer programming problem. The second is a FORTRAN program to solve the exchange heuristic algorithm discussed previously.

The Pennsylvania State University

The Graduate School

An Analysis of the Multiple Objective
Capital Budgeting Problem Via
Fuzzy Linear Integer (0-1) Programming

A Thesis in
Industrial Engineering and Operations Research

by

Michael G. Headly

Submitted in Partial Fulfillment
of the Requirements
for the Degree of

Master of Science

May 1980

Accession For	
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ABSTRACT

A multiple objective fuzzy linear programming approach to the capital budgeting problem is developed. Since much of the available data in any capital budgeting decision situation is either of an imprecise or ill-defined nature, a mathematical optimization technique is required that is capable of incorporating this inherent uncertainty. Fuzzy linear programming provides an effective methodology for this analysis.

Specifically, a mathematical model is developed which utilizes fuzzy linear programming as a solution technique for the research and development program or project selection problem. In addition, an exchange heuristic, a modified form of C. C. Petersen's exchange algorithm, is presented.

ACKNOWLEDGMENTS

The author expresses his gratitude to Dr. James Ignizio, his thesis adviser, for guidance during his work at The Pennsylvania State University.

The author is also grateful to Dr. Jack Hayya and Dr. Emory Ensore for their contribution to this thesis.

Support for the author's graduate work has been provided by the United States Army.

CHAPTER 1

INTRODUCTION

1.1 Purpose of the Research

The objective of the research documented in this thesis is the application and demonstration of a method for analysis of management decisions involving multiple objectives and constraints which are of a vague or ill-defined nature.

The traditional capital budgeting problem involves a single objective deterministic approach to the allocation of limited resources among available investment opportunities. The selection from among the various investment possibilities is such that the total return from the investment is maximized. In contrast to the traditional problem formulation, real-world capital investment decision analysis invariably encompasses nondeterministic systems involving multiple and usually conflicting objectives.

Investment selection or program selection in research and development planning is a multifaceted decision regularly faced by decision makers in government, industry, and the military. The constantly expanding nature of technological development necessitates decisions that involve multiple objectives in the decision criteria. Simply maximizing total return is an unrealistic and oversimplified decision criterion.

The complex selection process of research development programs may include the consideration of numerous factors, some of which are monetary while others are nonmonetary in nature. Influencing factors,

whose primary concern is not income generating, are demonstrated in safety and environmental considerations, which are inherent in virtually all business decisions today. The decision maker is clearly faced with a decision situation which is characterized mathematically as multiple criteria decision making.

Many mathematical programming techniques have been employed as a means of solving the capital budgeting problem; and, specifically, the investment or program selection problem has received a great deal of attention. A relatively new multiple objective optimization technique is fuzzy linear programming.

Fuzzy linear programming with its foundation in the theory of fuzzy sets is an optimization methodology designed for problems that are either too vague or too ill-defined to allow analysis by classical mathematical techniques. The inherent uncertainty which is ever present in any capital investment decision is the motivating influence in an examination of the applicability of fuzzy linear programming as a solution technique for the capital budgeting problem.

The design of this study encompasses five main objectives. These are:

1. Review various mathematical programming methodologies so as to establish applicability to the capital budgeting problem.
2. Evaluate the applicability of fuzzy linear programming as a solution technique for the capital budgeting problem.

3. Develop a fuzzy linear integer programming algorithm to solve the capital budgeting problem.
4. Apply the fuzzy linear integer programming algorithm to a representative problem.
5. Discuss extensions of this study and identify additional areas to which fuzzy programming techniques have applicability.

1.2 Organization of the Paper

This paper is organized as follows. Chapter 2 includes a historical perspective of various methodologies that have been employed as solution techniques for the capital budgeting problem. In Chapter 3, the basic elements of the theory of fuzzy sets are reviewed. Decision making in a fuzzy environment is discussed, and the model of fuzzy linear programming is presented in Chapter 4. In Chapter 5, the fuzzy capital budgeting model is presented along with the solution algorithm. Two example problems are solved. The results of the study are reviewed in Chapter 6, as well as possible extensions, and additional areas of applicability are suggested.

CHAPTER 2

HISTORICAL PERSPECTIVE OF THE CAPITAL BUDGETING PROBLEM

2.1 General

Decision makers have always sought a means of analyzing alternative investment possibilities in an efficient manner. The past twenty-five years have seen the development of analytical techniques to provide this analysis. The development of numerous quantitative analysis techniques has provided decision makers with a framework to more efficiently conduct this analysis. The usefulness of these quantitative techniques has been greatly extended with the ever-increasing accessibility of computers. While the computer's capability to analyze and store data has increased tremendously, the cost has steadily decreased. Today, the use of computer technology is widespread. Since the cost of many computers is no longer prohibitive, many small industries are utilizing quantitative analysis techniques that previously were reserved for government and large industries.

The classical approach to the analysis of alternate investment possibilities has been the maximization or minimization of a single objective function. Traditionally, this objective has been the maximization of profits or the minimization of costs. A significant amount of discussion has been generated concerning the classical approach and its inapplicability to today's complex decisions [1-11]. The basis of single objective function mathematical modelling is lost when it is recognized that real decision makers do not attempt to

optimize a single objective function. Rather, a solution is sought that satisfies the numerous objective functions that characterize a decision process. The solution is a compromise from among the various objective functions [5, 10, 12, 13]. The compromise is the result of the real-world limitations imposed on decision makers.

2.2 Survey of Related Literature

The multiple objective function optimization technique of fuzzy linear programming is a relatively new approach to multiple criteria decision making. Zimmermann [9, 10, 11] has shown the mathematical feasibility of this approach and its application to the media selection problem originally posed by Charnes et al. [14]. Two extensive bibliographies have been published on works related to fuzzy systems [15, 16]. A search of the literature failed to identify additional works dealing with the application of fuzzy linear programming as a multiple objective optimization technique. Kickert [17] has recently published a work detailing the various fuzzy theories and their impact on decision-making processes. Yager [18] discusses an eigenvector approach to the multiple objective optimization problem using fuzzy sets. There is increasing interest in multiple criteria decision making; and, correspondingly, a great deal of literature is available related to this work. The following paragraphs summarize a survey of the current literature on multiple criteria decision making, with an emphasis toward the capital budgeting problem.

A conference proceedings including numerous works on multiple criteria decision making was published by the University of South Carolina. A bibliography on multiple criteria decision making is included [19].

One mathematical programming technique that has been utilized for years as an optimization technique is linear programming. Charnes and Cooper [20] demonstrated an early use of linear programming as a solution technique for the problem of allocating funds. In recent years, multiple objective linear programming techniques have been developed. Benayoun, Larichev, de Montogolfier, and Tergny [21] discuss a methodology of using linear programming with multiple objectives. Belenson and Kapur [22] present an algorithm for solving multi-criteria linear programming problems with several examples. A multi-objective linear programming methodology has been presented by Evans and Steuer [1].

Goal programming is another robust optimization technique for dealing with decision problems involving multiple objectives. This technique was developed by Charnes and Cooper in the early 1950's [23]. Goal programming is an effective modelling methodology which affords an analysis of problems involving multiple, and possibly, conflicting objectives. The methodology requires an assignment of a priority to each objective. This priority assignment is a preemptive prioritization of the objectives in accordance with the priorities of the decision maker. Lee [13] published the first book entirely devoted to linear goal programming. Ijiri [24] in his work developed the concept of preemptive prioritization of objectives. Numerous applications of goal programming are available. These include capital budgeting optimization [4, 5, 8, 25, 26]; manpower planning [27]; academic planning, financial planning, and economic planning [13]; antenna array design and transportation problem [5]; and media

planning [14]. Survey works of goal programming have been published by Kornbluth [28] and Ignizio [3].

Integer and nonlinear goal programming algorithms have been developed and have realized many successful applications [5, 29]. Research is continuing to extend goal programming into the area of stochastic analysis. Contini [30] has demonstrated the mathematical feasibility of such an approach.

Interactive programming is yet another multi-criteria programming approach currently being utilized. The decision maker in this approach is required to specify trade-offs between the various objective functions. The process of specifying trade-offs is continued in a successive manner until no further trade-offs are desired by the decision maker. Geoffrion, Dyer, and Feinberg [31] demonstrate the application of interactive programming, while Zionts and Wallenius [32] present an overview of the interactive programming method as applied to the multiple criteria problem. Dyer [33] has also proposed an interactive goal programming technique, while Steuer [34, 35] has proposed an interactive approach to multiple objective linear programming.

Numerous other mathematical programming techniques have been discussed as solution methods for the multiple criteria decision problem. One technique that has received a great deal of attention is integer programming. The literature has many examples of the successful application of integer programming. Seward, Plane, and Hendrick [36] present an application in the area of allocating municipal funds for fire protection, Armstrong and Willis [37] discuss its use in the selection of water projects in California, and Nackel,

Goldman, and Fairman [38] demonstrate the use of integer programming in an example in the health care field. Chiu and Gear [39] present a stochastic integer programming approach to the research and development project selection problem.

A few of the other mathematical programming techniques with applications in the multiple criteria decision-making area are branch and bound procedures, dynamic programming and heuristic programming. Shih [40] has written on a branch and bound method, Kepler and Blackman [41] have demonstrated the use of dynamic programming in the selection of research and development projects, and Petersen [42, 43] has developed heuristic algorithms using exchange operations to solve the capital budgeting problem.

The recognition of the inherent risk and uncertainty in capital budgeting problems has been presented in many works in the literature. Hillier [44] presents a basic model for capital budgeting of risky interrelated projects. Stochastic analysis was initially proposed by Charnes and Cooper [45]. Their technique was termed chance-constrained programming. Healy [46] and Armstrong and Balintfy [47] have presented chance-constrained programming algorithms. Odom and Shannon [48] and Park and Theusen [49] have recently published works aimed at risk resolution in the capital budgeting decision analysis. Utility theory has also been a frequently employed technique in multicriteria decision making. Recent works in the literature include: Crawford, Huntzinger, and Kirkwood's [50] use of multiattribute utility theory in the selection of components of an electrical transmission system, and Keefer's [51] multiobjective analysis of research and development projects through the use of a multiattribute utility function.

The increased use of multiple criteria decision analysis is evident in the literature. Many excellent overviews of multiple objective optimization techniques are available. MacCrimmon [52] has analyzed the various techniques that are not mathematical programming approaches. These approaches involve either weighting factors methods, sequential elimination methods, or spatial proximity methods. Easton [2] reviews a variety of multivalued alternative weighting methods. Ignizio [3] reviews goal programming as a multiple objective optimization technique. Plane [53] presents integer programming and network analysis techniques, and Hax [54] discusses the use of decision analysis. Two survey papers [55, 56] discuss the use of the various decision-making techniques as related specifically to the capital budgeting problem.

The development of new approaches to the multiple criteria decision problem and the variety of applications of the more established techniques indicate a tremendous interest in multiple criteria optimization methodologies.

CHAPTER 3

BASIC FUZZY SET THEORY

3.1 The Decision-Making Process

The analysis of alternative courses of action culminating in a decision is an extremely complex process for the human mind. The complexity of real-world decision problems far exceed the capacity of the human mind to formulate and subsequently arrive at a reasonable solution [12]. The essence of a decision is that the decision maker is able to exercise his prerogative. Obviously, then, the decision maker must be faced with a situation involving several alternatives about which information is available. This information may be of a precise or exact type, or it may be vague or ill-defined. An effective decision-making process is normally an iterative process with a feedback capability so that, at various stages, additional information may enter into the analysis. The decision-making process with feedback is shown in Figure 1.

Decision making utilizing the multiple objective optimization technique of fuzzy linear programming is an effective methodology in which to employ this feedback process. Prior to any elaboration on fuzzy linear programming, a brief discussion of the basic principles of fuzzy set theory is necessary.

3.2 Fuzzy Set Theory

The theory of fuzzy sets was developed in response to a need for a conceptual framework to deal with problems which were either too

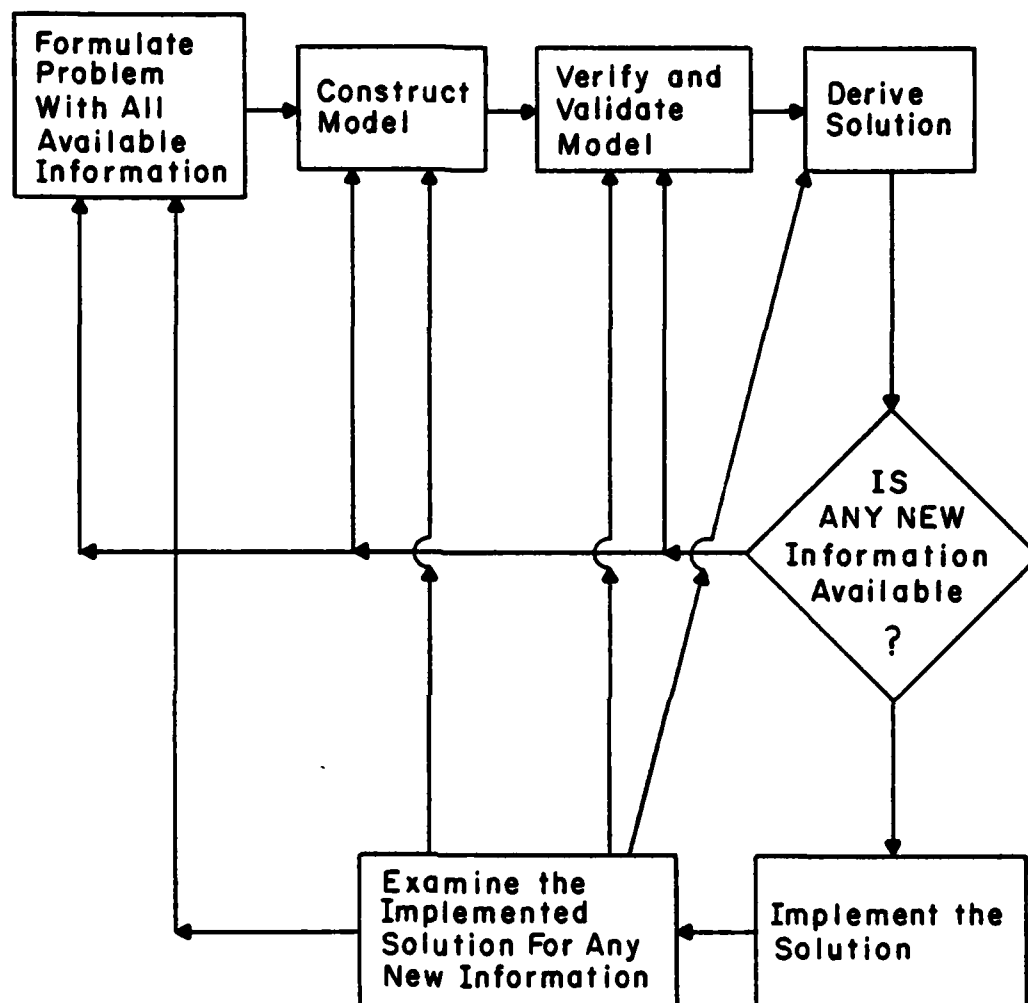


Figure 1. Decision Making as a Feedback Process

complex or too ill-defined to allow analysis by classical mathematical techniques.

Classical mathematics are much too rigid to be utilized in the optimization of systems that are humanistic in nature. These systems are composed largely of human perceptions and human judgments. Such systems are those in the fields of economics, psychology, sociology, linguistics, management science, medicine, law, philosophy, and others whose basic tenets are imprecise or fuzzy in nature.

The theory of fuzzy sets is founded on the theory of classes. Events may be viewed as in a continuum with respect to their membership or nonmembership in a class. The degree of membership in a class is the fundamental concept in the theory.

Classical mathematics' precise formulation of decision situations does not allow for the inclusion of a decision maker's judgmental capability. The concepts of fuzzy set theory create an overlap of the decision maker's judgmental ability and his quantitative analysis capabilities. The judgmental capability of the human mind analyzes a situation in an imprecise or approximate manner.

This imprecise or approximate analysis is necessitated by the complexity of today's managerial decision requirements. Real-life problems present themselves daily in vague or ill-defined ways. Many phenomena exist such as "satisfactory profits," "adequate return on investment," or "better productivity." None of these problems could be defined in precise mathematical terms. Instead, they would be twisted so as to conform to a precise mathematical optimization technique; and, therefore, the derived solution may or may not be accurate. In our attempts to understand and optimize systems which

are composed of various humanistic subsystems, the solutions obtained may pretend a higher degree of preciseness than is actually possible to achieve in the real system [57].

Fuzzy set theory provides a formal mathematical theory to analyze systems that are vague or inexact, with the vague or inexact nature defined by a fuzzy set [58].

3.3 Basic Definitions of the Theory of Fuzzy Sets

Zadeh [57] introduced the theory of fuzzy sets through the theory of sets, a generally universal mathematical theory. A set is defined as consisting of a finite or infinite number of elements [59]. The characteristic function of a set enables us to discuss the membership of the set in terms of functions. To define the characteristic function of a set, let A be a subset of the universe [60]. The function χ_A , the characteristic function, can only take on the values 0 or 1. If the universe is $X = \{x\}$, then χ_A is defined by the following:

$$\begin{aligned}\chi_A(x) &= 1 && \text{if } x \in A, \\ \chi_A(x) &= 0 && \text{if } x \notin A.\end{aligned}$$

Zadeh [57] utilized this concept of the characteristic function in his development of fuzzy set theory. Instead of the characteristic function being limited to only taking on the values 0 or 1, it is generalized to assume an infinite number of values between 0 and 1.

The basic definitions of fuzzy set theory which are important in the development of fuzzy linear programming will be presented in the following pages. These definitions are summarized from presentations by Zimmermann [9, 10, 11, 16] and Kickert [17].

Fuzzy Set - A class with a continuum of grades of membership.

Let X be a space of points (objects), with a generic element of X denoted by x_1 , then, $X = \{x\}$. The fuzzy set A in X is characterized by a membership function $\mu_A(x)$ which associated with each point in X a nonnegative real number whose supremum is finite, with $\mu_A(x)$ representing the grade of membership of A in X . This is represented as:

$$A = \{x, \mu_A(x) | x \in X\},$$

where $\mu_A(x)$ is the membership function of A in X .

Example: In the field of psychology, and specifically related to learning theory, the concepts of performance, learning, motivation, and anxiety are critical in the prediction of the outcome of any learning acquisition task.

Let

$$X = \{0, 10, 20, 30, 40, 50, 60, 70, 80, 90, 100\}$$

be possible scores which an individual may attain on a learning acquisition task. Fuzzy set A , "Motivation Levels Affecting Learning Acquisition," may be defined for a certain individual as:

$$A = \{(10, 0.2), (20, 0.4), (30, 0.6), (40, 0.65), (50, 0.7), (60, 0.75), (70, 0.85), (80, 1.0), (90, 0.9), (100, 0.8)\}.$$

Fuzzy set B , "Anxiety Levels Affecting Learning Acquisition," may be stated in a similar manner for the same individual as follows:

$$B = \{(10, 0.1), (20, 0.3), (30, 0.5), \\ (40, 0.60), (50, 0.65), (60, 0.75), \\ (70, 0.85), (80, 0.95), (90, 1.0), \\ (100, 0.85)\} .$$

Graphically, these two fuzzy sets are shown in Figure 2.

Intersection - In set theory, the intersection of two sets A and B, written $A \cap B$, is the set C containing all elements common to A and B. In fuzzy set theory, the membership function of $A \cap B$ is defined as:

$$\mu(x) = \text{Min} [\mu_A(x), \mu_B(x)] \quad \text{for all } x \in X .$$

Example: In the learning theory example, the fuzzy set representing the intersection of fuzzy sets A and B would be the fuzzy set C. Fuzzy set C is defined as:

$$C = \{(10, 0.1), (20, 0.3), (30, 0.5), \\ (40, 0.60), (50, 0.65), (60, 0.75), \\ (70, 0.85), (80, 0.95), (90, 0.90), \\ (100, 0.80)\} .$$

Union - In set theory, the union of two sets A and B, written $A \cup B$, is the set D containing all elements in either A or B, or both. In fuzzy set theory, the membership function of $A \cup B$ is defined as:

$$\mu(x) = \text{Max} [\mu_A(x), \mu_B(x)] \quad \text{for all } x \in X .$$

Example: In the learning theory example, the fuzzy set representing the union of fuzzy sets A and B would be the fuzzy set D. Fuzzy set D is defined as:

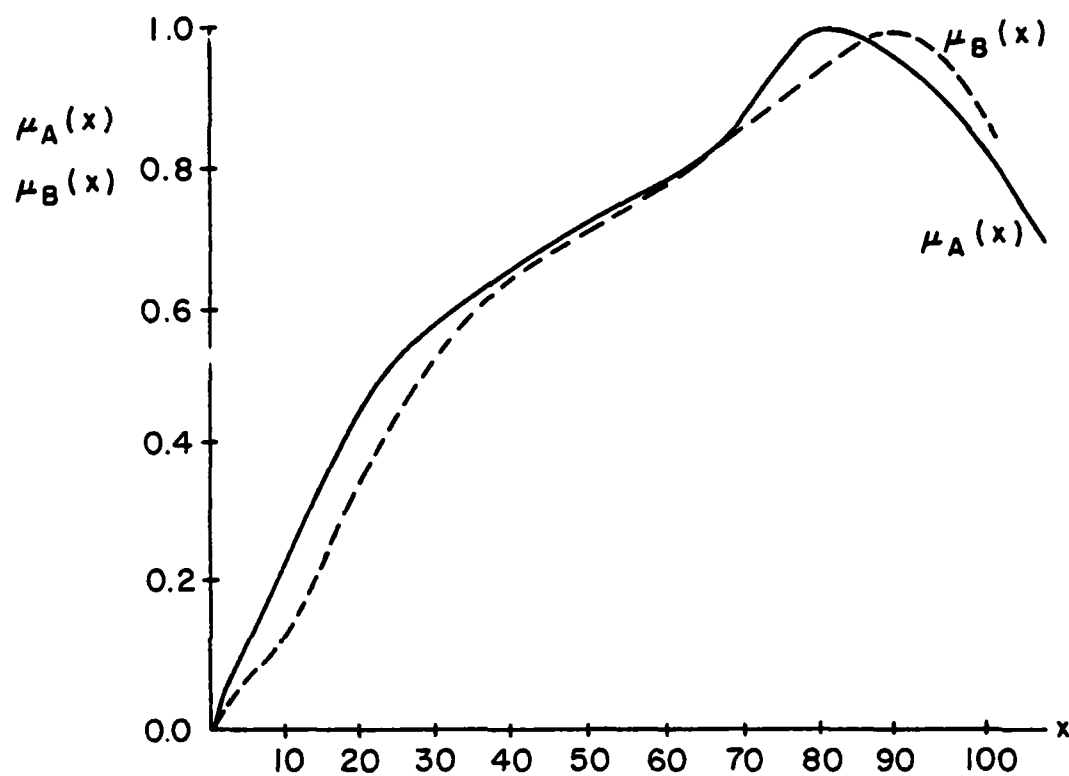


Figure 2. Illustration of Fuzzy Sets A and B

$$D = \{(10, 0.2), (20, 0.4), (30, 0.6), \\ (40, 0.65), (50, 0.7), (60, 0.75), \\ (70, 0.85), (80, 1.0), (90, 1.0), \\ (100, 0.85)\} .$$

The union of fuzzy sets A and B is displayed in Figure 3, and the intersection of the two fuzzy sets is shown in Figure 4.

Equality - Two fuzzy sets are equal if

$$\mu_A(x) = \mu_B(x) \quad \text{for all } x \in X .$$

Normality - The definition of the membership function did not limit the values $\mu(x)$ could assume. If the supremum of the membership function equals 1, then the fuzzy set is called normal. This is defined as:

$$\sup_X \mu_A(x) = 1 .$$

A fuzzy set can be normalized by dividing $\mu_A(x)$ by $\sup_X \mu_A(x)$.

Algebraic Product - The algebraic product of two fuzzy sets A and B is denoted AB and is defined in terms of the membership functions of the fuzzy sets A and B .

$$\mu_{AB}(x) = \mu_A(x) \cdot \mu_B(x) .$$

Algebraic Sum - The algebraic sum of two fuzzy sets A and B is denoted by $(A + B)$ and is defined in terms of the membership functions of the fuzzy sets A and B .

$$\mu_{(A+B)}(x) = \mu_A(x) + \mu_B(x) - \mu_A(x) \cdot \mu_B(x) .$$

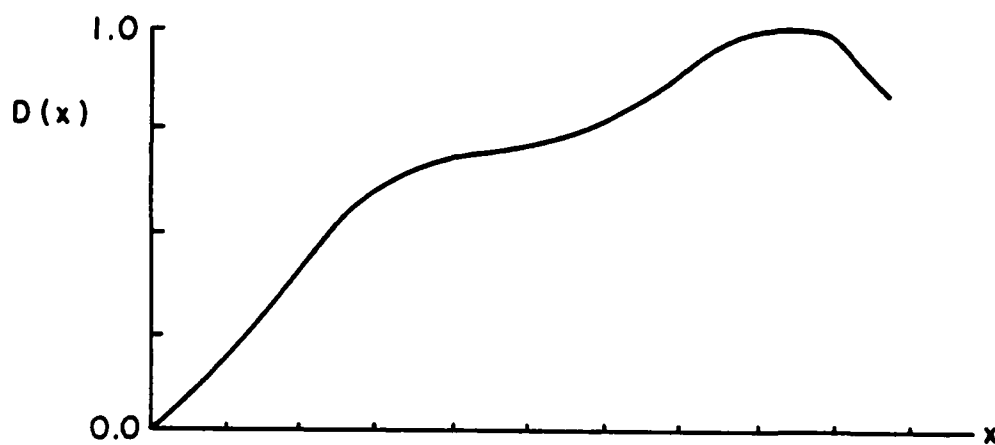


Figure 3. Union of Fuzzy Sets A and B

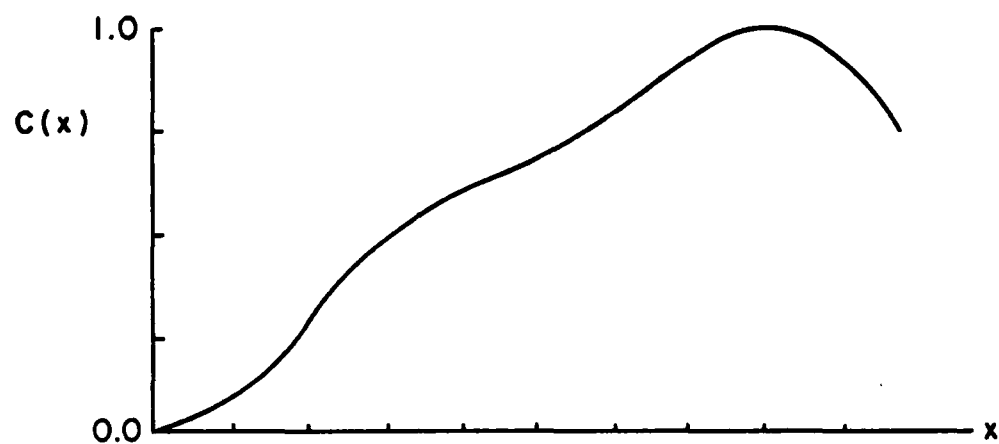


Figure 4. Intersection of Fuzzy Sets A and B

Containment - The fuzzy set definition of containment is analogous to the set theory definition of a subset. Fuzzy set A' is contained in fuzzy set B' if the membership function of A' is less than or equal to that of B' everywhere on X .

The basic definitions presented are sufficient for the discussion of fuzzy linear integer programming; however, there are many more concepts in the overall theory of fuzzy sets. For a more extensive treatment of the theory of fuzzy sets, Kaufmann [61] presents a complete review of the general theory of fuzzy sets.

CHAPTER 4

DECISION MAKING IN A FUZZY ENVIRONMENT

4.1 Fuzzy Decisions

In traditional decision making, the optimal decision is the selection of the activity or program with the highest desirability. In fuzzy decision making, the objective function(s) as well as the constraints may be fuzzy sets, each characterized by their membership functions. The optimal decision in the fuzzy environment is the fuzzy set formed by the intersections of the fuzzy sets describing the objective function(s) and constraints. Figure 5 illustrates the fuzzy decision process.

The region of intersection is a fuzzy set representing those activities which simultaneously satisfy the objective function(s) and the constraints. A solution to this fuzzy situation would be to select that point in the region of intersection with the greatest desirability or the highest degree of membership in the fuzzy set formed in the fuzzy decision. The selection of this solution point is analogous to the geometric representation of a solution to a linear programming problem [62]. The determination of the solution to the linear programming problem involving the intersection of n fuzzy sets is one of the basic principles in the development of fuzzy linear integer programming.

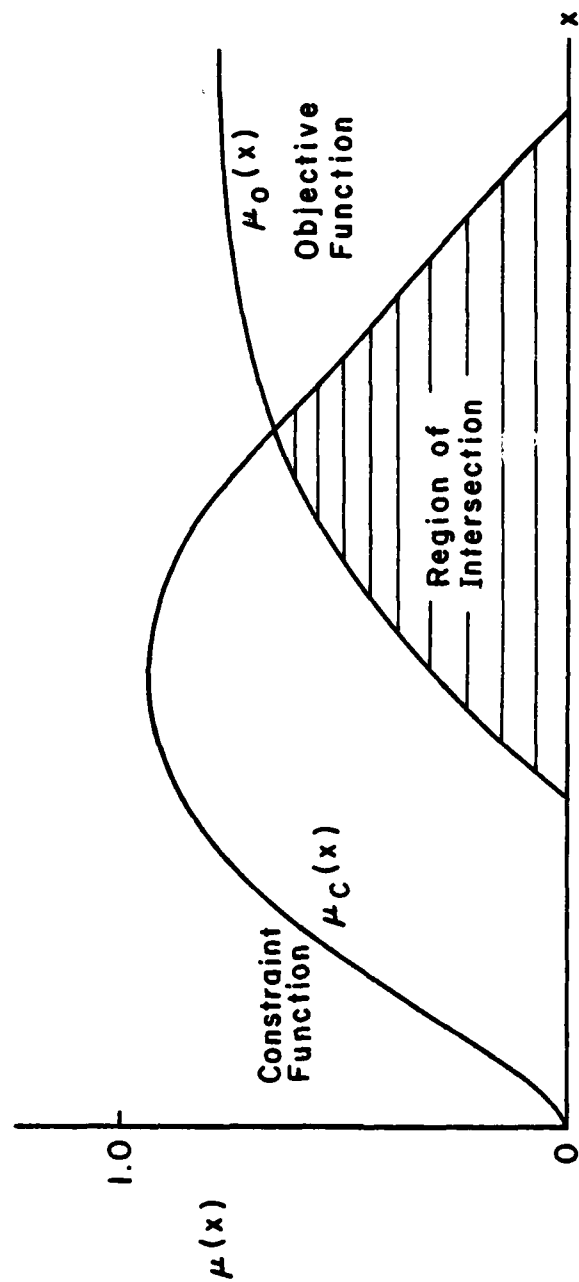


Figure 5. Fuzzy Decision Process

4.2 Fuzzy Linear Programming

The extension of fuzzy set theory into linear programming was utilized by Zimmermann [9]. The development of the fuzzy linear programming problem is as follows:

Start with the traditional vector minimization problem.

$$\text{Minimize } \bar{z} = \bar{c} \bar{x}$$

$$\text{Subject to } A\bar{x} \leq \bar{b}$$

$$\bar{x} \geq 0$$

The fuzzy version of this same linear programming problem is:

$$\bar{c} \bar{x} \leq \bar{z}^0$$

$$A \bar{x} \leq \bar{b}$$

$$\bar{x} \geq 0$$

where

\bar{c} is the vector of coefficients of the objective functions,

\bar{b} is the vector of constraints,

A is the coefficient matrix, and

\bar{z}^0 is the vector of aspiration levels of the fuzzy objectives and constraints.

The membership function $\mu(x)$ is defined such that it complies with the definition of a fuzzy set [57], that is, a real number in the interval (0,1).

$$\mu(x) = \begin{cases} 1 & \text{if } A\bar{x} \leq \bar{b} \text{ and } \bar{c} \bar{x} \leq \bar{z} \text{ is satisfied} \\ 0 & \text{if } A\bar{x} \leq \bar{b} \text{ and } \bar{c} \bar{x} \leq \bar{z} \text{ is strongly violated.} \end{cases}$$

The concept of an objective function being strongly violated or weakly violated is an important aspect of the decision-making process in a fuzzy environment. The membership function in Figure 6 will be utilized

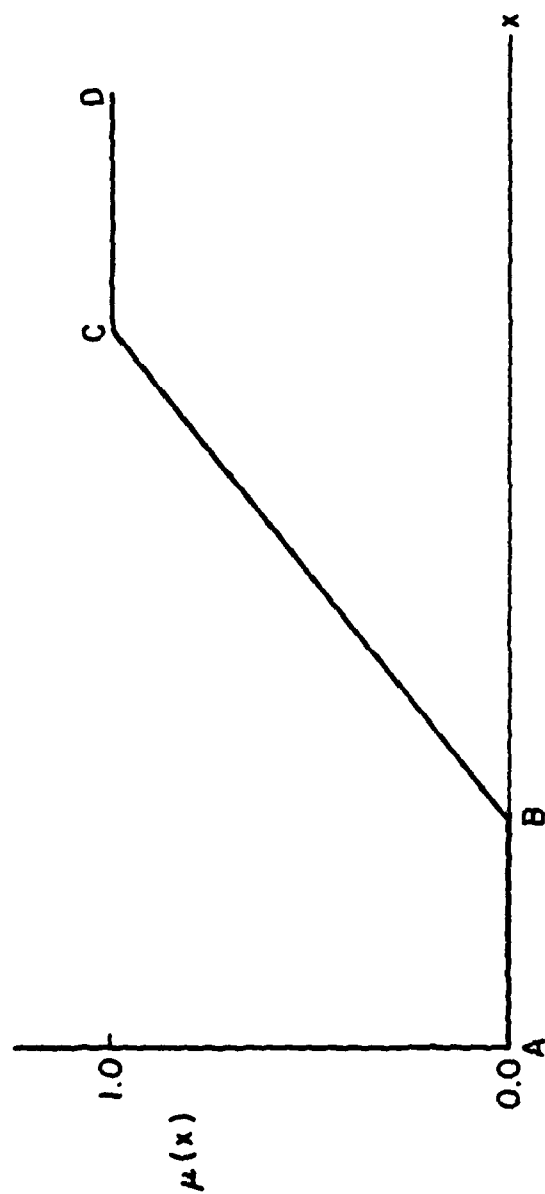


Figure 6. Membership Function of Fuzzy Set A

to illustrate this principle. Let this membership function be referred to as $\mu(x)$. In the interval CD, the membership function $\mu(x)$ is completely satisfied. The function describing the fuzzy set in this interval either achieves the aspiration level or exceeds it. In the interval BC, the membership function $\mu(x)$ is weakly violated. In this interval, the aspiration level is not achieved; however, the functional evaluation is greater than the lowest admissible value (Point B). The decision in this interval lies within the range of acceptable solutions as specified by the decision maker. In the interval AB, the membership function $\mu(x)$ is strongly violated. In this interval, any decision would lie wholly outside the acceptable range of solutions, since the functional evaluation of the fuzzy set would be less than the lowest admissible value, as specified by the decision maker.

If we let the fuzzy set B represent the intersection of the fuzzy sets representing the objective functions and the constraints, then the membership function of fuzzy set B is:

$$\mu_B(x) = \mu_O \cap \mu_C .$$

The intersection of these two fuzzy sets is defined by the min operator to be:

$$\mu(Bx) = \min_i \mu_i ; x \geq 0 .$$

The maximizing decision is simply

$$\max_{x \geq 0} \min_i [\mu_i(Bx)] ,$$

which minimizes the maximum violation of the membership function.

If the solution technique is to be linear programming, the following assumptions are necessary [11]:

1. All objective functions must have a specified aspiration level. The objective functions are expressed in the form

$$C_i \bar{x} \leq Z_i, \quad i = 1, 2, \dots, n.$$

2. If the objective functions are in the same form as the constraints, then the problem may be formulated in the following form:

$$A\bar{x} \leq \bar{b},$$

where

A is the matrix of coefficients, and

\bar{b} is the vector of aspiration levels of the objectives and the right-hand side values of the constraints.

3. The functions are assumed to be linear over the interval of consideration.

Given that assumption number (3) is satisfied, the linear membership function of fuzzy set B , the solution set of the intersection of the fuzzy sets representing the objectives and the constraints is:

$$\mu_B(x)_i = \begin{cases} 1 & \text{if } (Bx)_i \leq b'_i \\ 1 - \frac{(Bx)_i - b'_i}{d_i} & \text{if } b'_i < (Bx)_i \leq b'_i + d_i \\ 0 & \text{if } (Bx)_i > b'_i + d_i \end{cases},$$

where

- i indicates the i^{th} row of matrix B or b' ,
- B is A , the coefficient matrix, augmented by the rows of the objective functions,
- b' is the vector of the right-hand side values augmented by the upper bounds of the objective functions, and
- d_i is the subjectively selected value of admissible violation.

By substituting

$$b'_i = \frac{b'_i}{d_i} \quad \text{and} \quad B'_i = \frac{B_i}{d_i}$$

into the function $\mu_B(x)$, the maximizing decision then becomes:

$$\underset{x \geq 0}{\text{Max}} \quad \underset{i}{\text{Min}} [b'_i - (B'_i x)_i]$$

or

$$\underset{z \geq 0}{\text{Max}} \quad \mu_D(x),$$

where $\mu_D(x)$ represents the membership function of the fuzzy set representing the decision set.

It has been shown that the solution to this problem is equivalent to the following linear programming problem [9, 10, 11]:

$$\begin{aligned} &\text{Maximize} \quad \lambda \\ &\text{Subject to} \quad \lambda \leq b'_i - (B'_i x)_i, \quad i = 0, 1, \dots, n \\ &\quad \quad \quad x > 0. \end{aligned}$$

To demonstrate a continuous fuzzy linear programming problem, consider the following example:

$$\text{Maximize } Z = 4x_1 + 6x_2 + 8x_3 + 10x_4$$

$$\text{Subject to } x_1 + 3x_2 + 4x_3 + 2x_4 \leq 40$$

$$3x_1 + 2x_2 + 3x_3 + 6x_4 \leq 60$$

$$4x_1 + x_2 + 2x_3 + 3x_4 \leq 50$$

Solving this linear programming problem with the IBM MPSX mathematical programming system, the resulting program :

$$\bar{x} = (0, 8.57, 0, 7.14)$$

and

$$Z = 122.86$$

The problem when formulated into the fuzzy linear programming equivalent utilizing the subjectively selected d_i values follows. The aspiration levels and the lowest admissible values as well as the allowable admissible ranges are shown in Table IV.1.

Table IV.1. Selected Values for Fuzzy Transformations

	$\mu = 0$	$\mu = 1$	d_i
Objective function	115	140	25
First constraint	50	40	10

where

$\mu = 0$ decision maker specified lowest admissible value,

$\mu = 1$ decision maker specified aspiration level,

d_i decision maker specified range of acceptable values.

The resulting fuzzy linear programming formulation is:

$$\begin{aligned}
 &\text{Maximize } \lambda \\
 &\text{Subject to } \lambda \leq -4.6 + 0.16x_1 + 0.24x_2 + 0.32x_3 + 0.4x_4 \\
 &\quad \lambda \leq 5 - 0.1x_1 - 0.3x_2 - 0.4x_3 - 0.2x_4 \\
 &\quad 3x_1 + 2x_2 + 3x_3 + 6x_4 \leq 60 \\
 &\quad 4x_1 + x_2 + 2x_3 + 3x_4 \leq 50
 \end{aligned}$$

The solution to the fuzzy linear programming formulation is compared to the linear programming solution in Table IV.2. The fuzzy linear programming problem was solved using the IBM MPSX mathematical programming system.

Table IV.2. Summary of Calculations

<u>Linear Programming</u>	<u>Fuzzy Linear Programming</u>
$x_1 = 0.0$	$x_1 = 0.0$
$x_2 = 8.57$	$x_2 = 10.59$
$x_3 = 0.0$	$x_3 = 0.0$
$x_4 = 7.14$	$x_4 = 6.47$
$Z = 122.86$	$Z = 128.24$

The first advantage of fuzzy programming is that the decision maker is not required to specify in a precise manner the parameters of a decision situation. The decision maker is able to specify ranges of acceptability for those objective and constraint functions represented by fuzzy sets. In this example problem, the flexibility obtained in

the use of fuzzy linear programming enabled the decision maker to realize a greater return.

The second advantage of fuzzy programming is the ease with which it can be converted into a conventional mathematical programming problem. This is important due to the current availability of many mathematical programming techniques and algorithms [17].

CHAPTER 5

ZERO-ONE CAPITAL BUDGETING ALGORITHM

5.1 The Capital Budgeting Problem

The traditional capital budgeting problem involves a single objective function deterministic approach to the allocation of limited resources among available investment opportunities. This approach differs greatly from most real-world capital budgeting problems. Actual resource allocation distribution procedures involve an analysis which is by necessity nondeterministic and sensitive to numerous conflicting interests. Due in part to this divergence between the traditional mathematical model of the capital budgeting problem and the necessities of real-world decision making, a significant amount of discussion has been generated concerning the traditional approach and its applicability to today's complex decision-making procedures [1-11].

The solution to the capital budgeting problem obtained in a model which seeks a compromise from among the numerous objective functions which represent the decision situation is a more viable methodology to characterize today's complex decision-making situations [5, 10, 12, 13]. Rather than a single objective function model of the capital budgeting problem, the general multiple objective function model takes on the following form:

$$\text{Maximize} \quad \sum_{j=1}^n r_{kj} x_j \quad k = 1, 2, \dots, K$$

$$\text{Subject to} \quad \sum_{j=1}^n c_{ij} x_j \leq b_i \quad \forall i$$

$$x_j = (0,1) \quad ,$$

where the terms are defined as:

$$x_j = \begin{cases} 1 & \text{if the } j^{\text{th}} \text{ alternative is selected} \\ 0 & \text{if the } j^{\text{th}} \text{ alternative is not selected,} \end{cases}$$

$$r_{kj} = \text{return on objective } k \text{ from alternative } j \text{ ,}$$

$$c_{ij} = \text{requirement of resource } i \text{ by alternative } j \text{ , and}$$

$$b_i = \text{limitation of resource } i \text{ .}$$

Many multiple objective optimization techniques have been employed in the solution of this problem; these were discussed in Chapter 2.

5.2 Fuzzy Linear Integer Programming/Exchange Heuristic Algorithm

An algorithm is developed which combines the principles of fuzzy linear programming and Petersen's [42] exchange heuristic to solve the multiple objective capital budgeting problem. The algorithm is intended to solve the following capital budgeting problem:

$$\text{Maximize} \quad \sum_{j=1}^n r_{kj} x_j \quad k = 1, 2, \dots, K$$

$$\text{Subject to} \quad \sum_{j=1}^n c_{ij} x_j \leq b_i \quad \forall i$$

$$x_j = (0,1) \quad ,$$

where the terms are defined previously.

The algorithm is a three-phase solution technique which incorporates an interactive process between the analyst and the decision maker in Phase I. In Phase II, a fuzzy linear integer problem is solved. Phase III, the exchange heuristic, is utilized if a 0,1 solution was not obtained in Phase II.

5.3 The Algorithm

5.3.1 Phase I: Determination of aspiration levels and the lowest admissible values. Phase I of the algorithm is intended to be an interactive process between the analyst and the decision maker. In this phase, K successive linear programming problems are solved, where K is the number of fuzzy objectives. The constraint set is to remain constant throughout the evaluations. In this manner, each objective function yields the highest attainable value possible. This value will be referred to as the Aspiration Level.

The lowest admissible value for each function is determined from the programs which yield the aspiration levels for the other $K-1$ functions. The value determined to be the lowest admissible value when subtracted from the aspiration level yields the allowable tolerance interval for each objective function.

The calculated values for the aspiration levels, lowest admissible values, and the tolerance intervals should then be reviewed by the decision maker. It rests with the decision maker to provide the analyst with the values to continue the algorithm in Phase II. This interactive process is critical to the fundamental concept of fuzzy programming, that the theory of fuzzy sets combines the quantitative aspects of optimization with the judgmental abilities of decision makers.

The programming procedure utilized to complete this phase is the IBM MPSX Linear Programming technique. Appendix A discusses the IBM MPSX system in greater detail.

5.3.2 Phase II: Determination of a fuzzy linear integer programming solution. In Phase II, a fuzzy transformation is carried out on each fuzzy function, and a linear integer programming problem is solved to maximize the value of the membership function.

The fuzzy transformation depends on the type of function under consideration. The three possibilities are shown in Table V.1. The \bar{d}_i and d_i are the selected upper and lower bounds of the tolerance interval specified by the decision maker. Graphically, these three functions are shown in Figures 7, 8, and 9.

Table V.1. Fuzzy Transformations

<u>Type</u>	<u>Objective</u>	
I.	Equal or exceed b_i	$\lambda \leq 1 - \frac{b'_i - (Zx)_i}{d_i}$
II.	Equal or less than b_i	$\lambda \leq 1 - \frac{(Zx)_i - b'_i}{\bar{d}_i}$
III.	Equal b_i	a. $\lambda \leq 1 - \frac{b'_i - (Zx)_i}{d_i}$ and b. $\lambda \leq 1 - \frac{(Zx)_i - b'_i}{\bar{d}_i}$

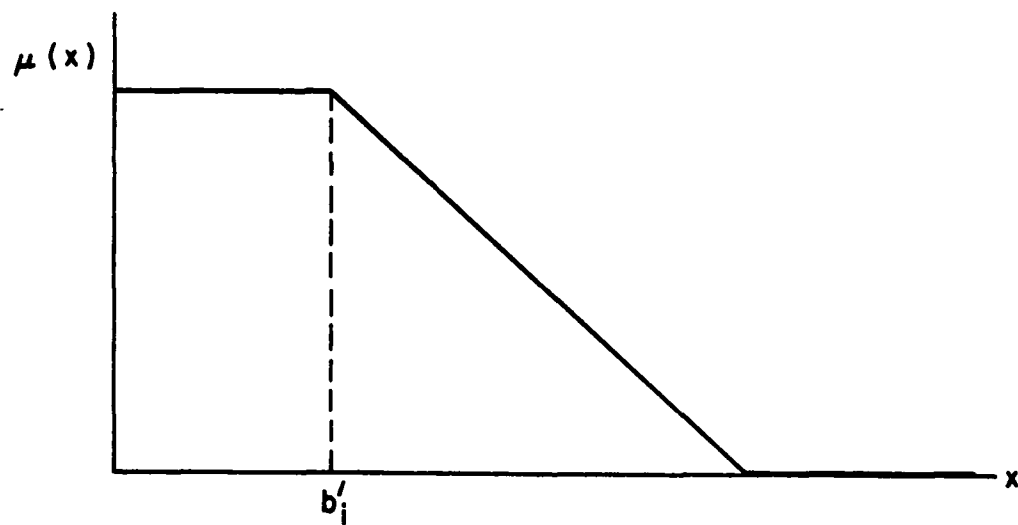


Figure 7. Membership Function for Type I Objective

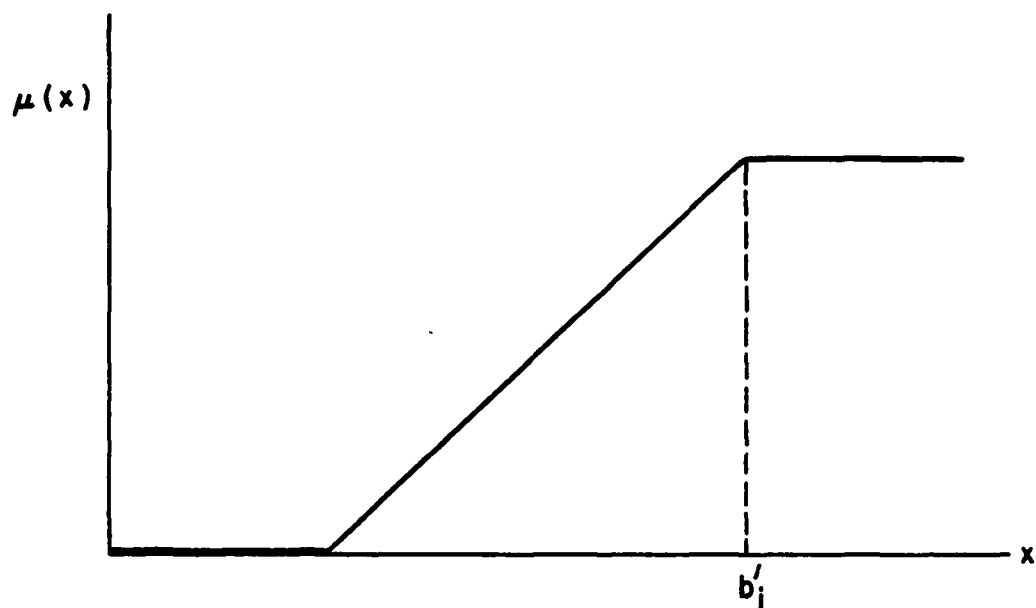


Figure 8. Membership Function for Type II Objective

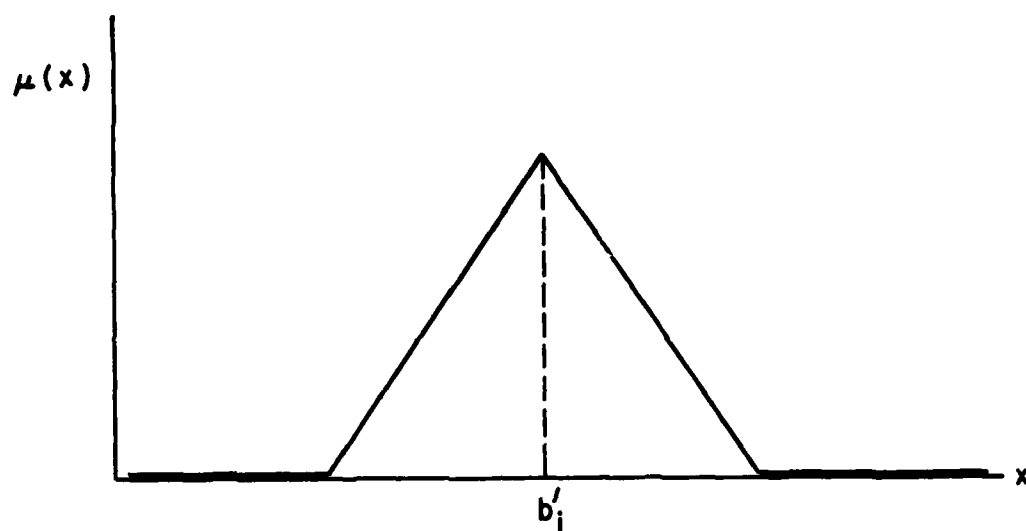


Figure 9. Membership Function for Type III Objective

The fuzzy linear integer programming problem formulation typically may be expressed as follows:

Maximize λ

Subject to $\lambda \leq 1 - \frac{b'_1 - (Zx)_1}{d_1}$, $i = 1, \dots, K$

and $(Ax)_1 \leq b_1$,

$x_j \leq 1 \quad \forall j$,

where $(Ax)_1$ is the set of rigid constraint functions, and each variable has an upper bound of 1.0.

If the solution to this linear programming problem is satisfactory to satisfy the 0-1 restrictions, then the algorithm terminates; otherwise, proceed to Phase III.

5.3.3 Phase III: Determination of an exchange heuristic 0-1 solution. The exchange heuristic, a modified form of Petersen's [42], is composed of three major steps:

1. Determination of an initial solution.
- ii. Determination of a fitback solution.
- iii. Utilize exchange operations progressively to improve the solution so as to finally achieve at least a local optimum.

Determination of an Initial Solution - The initial solution is obtained after ranking each variable based on the value of the ratio T_j/R_j given n variables and m objective functions, where T_j is the summation of the coefficient

values of each variable in the fuzzy objectives, and

R_j is defined as $\sum_{i=1}^n (c_{ij}/b_i)$ $j = 1, 2, \dots, n$.

The variables with the highest values of the ratio T_j/R_j are placed at the top of the ranking list.

Variables are rejected from the bottom of the list until

the rejection of the K^{th} variable causes satisfaction of

$\sum_{j=1}^{K-1} c_{ij} \leq b_i$ for all rigid constraints. The initial

solution is comprised of those variables ranked 1 through

$K-1$.

Determination of a Fitback Solution - In general,

following the selection of an initial solution, there

will be some degree of slack for each constraint. The

fitback solution selects from the initially nonselected

variables ranked $K+1$ to n , one or more that can be

included with the selected variables without violating

any constraint.

Exchange Operations - The alternatives in the sets of

selected and nonselected variables are ranked according

to their T_j value. In the set of selected variables,

the variables are ranked starting with the lowest value

first, while in the set of nonselected variables, variables

are ranked with the highest value first.

The search procedure is a two-step process. For

each exchange, it is determined if the exchange under

consideration would cause an improvement in the membership

function. If an improvement is noted, then the feasibility of the exchange is examined.

The set of exchanges is divided into two groups. The first search consists of the 2/1, 1/1, and 1/2 exchanges, while the second search considers 3/1, 3/2, and 3/3 exchanges. In each case, the first number refers to the number of variables selected from the set of nonselected variables.

The sequencing of the variables in the sets of selected and nonselected variables is performed to reduce the number of searches necessary to obtain a solution. The sequence allows for the examination of the most profitable exchanges first. Then, if an exchange is advantageous, the search is reduced due to dominance. In ordering the sets of selected and nonselected variables, the search proceeds naturally from the most advantageous exchanges to least advantageous exchanges.

5.4 Example of Three-Phase Algorithm

The three-phase algorithm is most easily explained via an example. Consider a problem in which the decision-making situation is characterized by two fuzzy objective functions and three rigid constraint functions. Assume this decision has the following problem formulation:

$$\text{Maximize } Z_1 = 3x_1 + 5x_2 + 5x_3 + x_4$$

$$\text{Maximize } Z_2 = x_1 + x_3 + x_4$$

$$\text{Subject to } 2x_1 + x_2 + 3x_3 + x_4 \leq 6$$

$$x_1 + 2x_2 + 4x_3 + 2x_4 \leq 5$$

$$3x_1 + 2x_2 + x_3 + x_4 \leq 4$$

$$x_j = (0,1) \quad .$$

5.4.1 Determine the Aspiration Level and Lowest Admissible Value for each objective. To calculate the aspiration level of the objective functions, the optimization technique of linear programming is utilized. Solving a linear programming problem to maximize each objective function subject to the same set of constraint functions yields the highest attainable value of the solution or the aspiration level. Thus, for the example:

$$(a) \text{ Maximize } Z = 3x_1 + 5x_2 + x_4$$

$$\text{Subject to } 2x_1 + x_2 + 3x_3 + x_4 \leq 6$$

$$x_1 + 2x_2 + 4x_3 + 2x_4 \leq 5$$

$$3x_1 + 2x_2 + x_3 + x_4 \leq 4$$

$$\text{Solution: } Z_1 = 9.55$$

$$x_1 = 0.45 \quad x_3 = 0.64$$

$$x_2 = 1.00 \quad x_4 = 0.0 \quad .$$

$$\begin{aligned}
 \text{(b) Maximize } Z_2 &= x_1 + x_3 + x_4 \\
 \text{Subject to } 2x_1 + x_2 + 3x_3 + x_4 &\leq 6 \\
 x_1 + 2x_2 + 4x_3 + 2x_4 &\leq 5 \\
 3x_1 + 2x_2 + x_3 + x_4 &\leq 4
 \end{aligned}$$

$$\begin{aligned}
 \text{Solution: } Z_2 &= 5.36 \\
 x_1 &= 0.82 & x_3 &= 0.55 \\
 x_2 &= 0.0 & x_4 &= 1.00
 \end{aligned}$$

To calculate the lowest admissible value for each of n objective functions, evaluate each objective function with the other $n - 1$ linear programming solution programs. Select as the lowest admissible value for each objective function the minimum resulting evaluation.

Thus, for the example:

- (a) Evaluate objective function Z_1 with the program obtained in Item (b) of the determination of the aspiration level.

$$\begin{aligned}
 Z_1 | (0.82, 0, 0.55, 1.0) &= 3x_1 + 5x_2 + 5x_3 + x_4 \\
 Z_1(\text{LAV}) &= 6.21
 \end{aligned}$$

- (b) Evaluate objective function Z_2 with the program obtained in Item (a) of the aspiration level.

$$\begin{aligned}
 Z_2 | (0.45, 1.0, 0.64, 0) &= x_1 + x_3 + x_4 \\
 Z_2(\text{LAV}) &= 1.09
 \end{aligned}$$

The results of Phase I of the algorithm are summarized in Table V.2.

Table V.2. Summary of Calculations in Phase I

<u>Objective Function</u>	<u>Aspiration Level</u>	<u>Lowest Admissible Value</u>	<u>Tolerance Interval</u>
Z_1	9.55	6.21	3.34
Z_2	5.36	1.09	4.27

5.4.2 Phase II: Fuzzy linear integer programming formulation.

The initial step in the fuzzy linear integer programming problem formulation is to determine the type of objective function and transform the objective function as appropriate. The fuzzy transformations were shown in Table V.1. Thus, for the example:

Since both fuzzy functions are Type I functions, the transformations are as follows:

$$Z_1: \lambda \leq 1 - \frac{[9.55 - (3x_1 + 5x_2 + 5x_3 + x_4)]}{3.34}$$

$$\lambda \leq 1 - (2.85 - 0.8982x_1 - 1.497x_2 - 1.497x_3 - 0.299x_4)$$

$$\lambda \leq -1.85 + 0.8982x_1 + 1.497x_2 + 1.497x_3 + 0.299x_4$$

$$Z_2: \lambda \leq 1 - \frac{[5.36 - (x_1 + x_3 + 4x_4)]}{4.27}$$

$$\lambda \leq 1 - (1.255 - 0.2342x_1 - 0.2342x_3 - 0.9367x_4)$$

$$\lambda \leq -0.255 + 0.2342x_1 + 0.2342x_3 + 0.9367x_4$$

The formulation as a fuzzy linear integer programming problem is:

$$\begin{aligned}
&\text{Maximize} && \lambda \\
&\text{Subject to} && \lambda \leq -1.85 + 0.8982x_1 + 1.497x_2 + 1.497x_3 + 0.299x_4 \\
&&& \lambda \leq -0.255 + 0.2342x_1 + 0.2342x_3 + 0.9367x_4 \\
&&& 2x_1 + x_2 + 3x_3 + x_4 \leq 6 \\
&&& x_1 + 2x_2 + 4x_3 + 2x_4 \leq 5 \\
&&& 3x_1 + 2x_2 + x_3 + x_4 \leq 4 \\
&&& x_j = (0,1) \quad .
\end{aligned}$$

$$\text{Solution: } \lambda = 0.21$$

$$\begin{aligned}
x_1 &= 1.0 & x_3 &= 1.0 \\
x_2 &= 0.0 & x_4 &= 0.0 \quad .
\end{aligned}$$

The solution is in the form (0,1) ; however, Phase III will be utilized to illustrate the exchange heuristic.

5.4.3 Phase III: Exchange Heuristic solution. The first step in this Exchange Heuristic approach is to set up the problem in the standard form as described previously (Section 5.2). In the example under consideration, this formulation is as follows:

$$\begin{aligned}
&\text{Maximize} && \lambda \\
&\text{Subject to} && 2x_1 + x_2 + 3x_3 + x_4 \leq 6 \\
&&& x_1 + 4x_3 + 4x_3 + 2x_4 \leq 5 \\
&&& 3x_1 + 2x_2 + x_3 + x_4 \leq 4 \\
&&& \lambda - 0.8982x_1 - 1.497x_2 - 1.497x_3 - 0.299x_4 \leq -1.85 \\
&&& \lambda - 0.2342x_1 - 0.2342x_3 - 0.9367x_4 \leq -0.255 \\
&&& x_j = (0,1) \quad .
\end{aligned}$$

The second step is to determine an initial solution. The variables comprising the initial solution are determined by ranking each variable based on the ratio T_j/R_j . T_j is defined as the summation of the coefficients of each variable in the fuzzy objective functions. In the example, there are two fuzzy objectives. Thus, in the example, the calculation of the T_j values is as follows:

Table V.3. Calculation of T_j Values

Variable	Fuzzy Objective 1	Fuzzy Objective 2	$T_j = \sum c_{1j}$
x_1	-0.8982	-0.2342	-1.1324
x_2	-1.497	0.0	-1.497
x_3	-1.497	-0.2342	-1.7312
x_4	-0.2990	-0.9367	-1.2357

The R_j values are calculated by evaluating the ratio of the coefficients of both the fuzzy objective functions and the rigid constraint functions to the appropriate b_i values. In the example problem, the calculation of the R_j values is as follows:

Table V.4. Calculation of R_j Values

<u>Variable</u>	<u>R_j</u>	<u>R_j Values</u>
x_1	$\frac{-0.8982}{-1.85} + \frac{-0.2342}{-0.255} + \frac{1}{5} + \frac{3}{4} + \frac{2}{6}$	2.69
x_2	$\frac{-1.497}{-1.85} + 0 + \frac{2}{5} + \frac{2}{4} + \frac{1}{6}$	1.876
x_3	$\frac{-1.497}{-1.85} + \frac{-0.2342}{-0.255} + \frac{4}{5} + \frac{1}{4} + \frac{3}{6}$	3.27
x_4	$\frac{-0.2990}{-1.85} + \frac{-0.9367}{-0.255} + \frac{2}{5} + \frac{1}{4} + \frac{1}{6}$	4.65

The ratio T_j/R_j is calculated from the results obtained in Table V.3 and Table V.4. In the example problem, the T_j/R_j values are:

Table V.5. Calculation of T_j/R_j Values

<u>Variable</u>	<u>T_j</u>	<u>R_j</u>	<u>T_j/R_j</u>
x_1	-1.1324	2.69	-0.420
x_2	-1.497	1.876	-0.798
x_3	-1.7339	3.27	-0.530
x_4	-1.2357	4.65	-0.266

The initial solution may be calculated by ranking the n variables according to their respective T_j/R_j values. Variables with the highest values of the T_j/R_j ratio are placed at the top of the ranking list. Variables are rejected from the bottom of the list until the rejection of the K^{th} variable causes satisfaction of $\sum_{j=1}^{k-1} c_{ij} \leq b_i$

for all rigid constraints. The initial solution is comprised of those variables ranked 1 through $K-1$. The initial ranking of the variables is as follows:

Table V.6. Initial Ranking of Variables

<u>Variable</u>	<u>Initial Ranking</u>
x_1	2
x_2	4
x_3	3
x_4	1

The initial solution may be determined as follows:

Table V.7. Determination of the Initial Solution

<u>Objective Function</u>	<u>Functional Evaluation</u>				
	<u>{4,1,3,2}</u>	<u>{4,1,3}</u>	<u>{4,1}</u>		
1	7*	6	3	Initial Solution	{4,1}
2	9*	7*	3	Set of Selected Variables:	{4,1}
3	7*	5*	4	Set of Nonselected Variables:	{2,3}
4	-4.19	-0.269	-1.19	Value of Membership Function:	0.0
5	-1.41	-1.40	-1.17		

* Indicates constraint functions that are not satisfied.

The third step in the exchange heuristic is to determine a fitback solution. The fitback solution selects from the set of initially nonselected variables one or more than can be included with the selected variables without violating any constraint. Table V.8 displays the calculation of the fitback solution.

Table V.8. Determination of Fitback Solution

Objective Function	<u>Functional Evaluation</u>		
	<u>{4,1,3}</u>	<u>{4,1,2}</u>	
1	6	4	Since at least one constraint is violated in each possible fitback solution, therefore, the fitback solution is the same as the initial solution, {4,1}.
2	7*	5	
3	5*	6*	
4	-2.69	-2.69	
5	-1.40	-1.40	

* Indicates constraint functions that are not satisfied.

The fourth step consists of utilizing exchange operations progressively to improve the solution. After each exchange, the feasibility of the exchange is examined. If the exchange is feasible, the possible improvement in the membership function is examined. If an improvement is not noted, then proceed to the next exchange possibility. Table V.9. displays the search procedure examining the possible exchanges.

Table V.9. First Search Exchange Procedure

<u>First Search</u>				
List of Selected Variables: {4,1}		List of Nonselected Variables: {2,3}		
<u>Attempted Exchange</u>		<u>λ</u>	<u>λ max</u>	<u>Selected Variables</u>
<u>(2/1)</u>			0.0	{4,1}
(2,3) for 4	Infeasible Exchange	--	0.0	{4,1}
(2,3) for 1	Infeasible Exchange	--	0.0	{4,1}
<u>(1/1)</u>				
(2) for 4	Infeasible Exchange	--	0.0	{4,1}
(2) for 1	Infeasible Exchange	--	0.0	{4,1}
(3) for 4	Advantageous/ Feasible	0.21	0.21	{3,1}

Repeat First Search

List of Selected Variables: {3,1}

List of Nonselected Variables: {2,4}

The algorithm utilizes a second pass through the list of non-selected variables once a favorable exchange has been noted. In the example problem, the second search produced no exchange advantageous to the membership function. Therefore, the final solution to the example is:

Set of Selected Variables: {3,1}

Value of Membership Function: $\lambda = 0.21$

5.5 Example Number 2, Project Selection Example

In order to illustrate a capital budgeting problem in which the decision situation is program or project selection, the following example is presented. In this example, the decision maker has specified firm values for the aspiration levels and acceptable ranges of admissibility for each fuzzy function [6].

A systems engineer has to design an integrated system composed of three subsystems, designated A, B, and C. Three systems have been proposed for Subsystem A, four for Subsystem B, and three for Subsystem C. Four attributes were established by management to guide in the selection of the subsystems. These are weight, development costs, estimated reliability, and power requirements. Table V.10 summarizes the attribute characteristics for each proposed candidate.

Design incompatibilities exist between Candidates A-2 and C-9. Also, due to design features, if Candidate B-7 is selected, then Candidate C-10 must be selected.

Management has established firm values for the aspiration levels and allowable ranges of admissibility. Table V.11 summarizes this information for each attribute.

Table V.10. Attribute Data for Proposed Candidates

<u>Attribute</u>	<u>Subsystem A</u>			<u>Subsystem B</u>			<u>Subsystem C</u>			
	<u>1</u>	<u>2</u>	<u>3</u>	<u>4</u>	<u>5</u>	<u>6</u>	<u>7</u>	<u>8</u>	<u>9</u>	<u>10</u>
Weight (lb.)	32	57	19	95	107	61	48	23	10	15
Cost (\$10 ⁴)	120	95	160	64	67	96	119	42	36	70
Reliability	0.97	0.94	0.99	0.89	0.90	0.94	0.96	0.98	0.97	0.99
Power (watts)	21	35	10	60	83	27	50	12	7	16

Table V.11. Management Specified Attribute Data

<u>Attribute</u>	<u>Aspiration Level</u>	<u>Lowest Admissible Value</u>	<u>Highest Admissible Value</u>
Weight (lb.)	150	120	165
Cost (\$10 ⁴)	195	260	- -
Power (watts)	100	70	110

Mathematically, this problem may be formulated as a system of linear equations. The reliability constraint may be transformed to a linear equation via the transformation $Z = (Y)^x = Z \ln Y$:

$$Z_1 = 32x_1 + 57x_2 + 19x_3 + 95x_4 + 107x_5 + 61x_6 + 48x_7 \\ + 23x_8 + 10x_9 + 15x_{10}$$

$$Z_2 = 120x_1 + 95x_2 + 160x_3 + 64x_4 + 67x_5 + 96x_6 + 119x_7 \\ + 42x_8 + 36x_9 + 70x_{10}$$

$$Z_3 = 21x_1 + 35x_2 + 10x_3 + 60x_4 + 83x_5 + 27x_6 + 50x_7 \\ + 12x_8 + 7x_9 + 16x_{10}$$

Subject to

$$(0.97)^{x_1} (0.94)^{x_2} (0.99)^{x_3} (0.89)^{x_4} (0.90)^{x_5} (0.94)^{x_6}$$

$$(0.96)^{x_7} (0.98)^{x_8} (0.97)^{x_9} (0.99)^{x_{10}} \geq 0.85$$

$$x_1 + x_2 + x_3 = 1$$

$$x_4 + x_5 + x_6 + x_7 = 1$$

$$x_8 + x_9 + x_{10} = 1$$

$$x_2 + x_5 = 1$$

$$x_7 - x_{10} = 0$$

$$x_j = (0,1)$$

Transforming the problem into its fuzzy equivalent, the problem formulation is as follows:

Maximize λ

Subject to

$$\lambda + 1.84x_1 + 1.46x_2 + 2.46x_3 + 0.98x_4 + 1.03x_5 + 1.47x_6 \\ 1.83x_7 + 0.64x_8 + 0.55x_9 + 1.08x_{10} \leq 4.0$$

$$\lambda - 1.067x_1 - 1.9x_2 - 0.64x_3 - 3.17x_4 - 3.56x_5 - 2.03x_6 \\ - 1.6x_7 - 0.76x_8 - 0.33x_9 - 0.5x_{10} \leq -4.0$$

$$\lambda + 2.13x_1 + 3.8x_2 + 1.26x_3 + 6.34x_4 + 7.13x_5 + 4.06x_6 \\ + 3.2x_7 + 1.53x_8 + 0.67x_9 + 1.0x_{10} \leq 11.0$$

$$\lambda - 0.70x_1 - 1.16x_2 - 0.34x_3 - 2.0x_4 - 2.77x_5 - 0.9x_6 \\ - 1.66x_7 - 0.4x_8 - 0.23x_9 - 0.53x_{10} \leq -2.33$$

$$\lambda + 2.1x_1 + 3.5x_2 + 1.0x_3 + 6.0x_4 + 8.3x_5 + 2.7x_6 + 5.0x_7 \\ + 1.2x_8 + 0.7x_9 + 1.6x_{10} \leq 11.0$$

$$\lambda - 0.032x_1 - 0.062x_2 - 0.01x_3 - 0.117x_4 - 0.105x_5 - 0.062x_6 \\ - 0.041x_7 - 0.020x_8 - 0.030x_9 - 0.01x_{10} \geq -0.1625$$

$$1.0x_2 + 1.0x_5 = 1.0$$

$$1.0x_1 + 1.0x_2 + 1.0x_3 = 1.0$$

$$1.0x_4 + 1.0x_5 + 1.0x_6 + 1.0x_7 = 1.0$$

$$1.0x_8 + 1.0x_9 + 1.0x_{10} = 1.0$$

$$1.0x_7 - 1.0x_{10} = 0$$

$$x_j = (0,1)$$

The solution to the fuzzy linear programming problem is:

$$\lambda = 0.1344$$

$$x_1 = 0.0$$

$$x_6 = 1.0$$

$$x_2 = 1.0$$

$$x_7 = 0.0$$

$$x_3 = 0.0$$

$$x_8 = 1.0$$

$$x_4 = 0.0$$

$$x_9 = 0.0$$

$$x_5 = 0.0$$

$$x_{10} = 0.0$$

The solution is in the required 0-1 form, the exchange heuristic or Phase III is not necessary. Therefore the decision is summarized as follows:

$$\lambda = 0.1334$$

List of Selected Projects: {2,6,8}

List of Nonselected Projects: {1,3,4,5,7,9,10} .

The rigid constraints were all satisfied. The fuzzy objectives were calculated as follows:

Table V.12. Example 2, Attribute Satisfaction

<u>Attribute</u>	<u>Calculated Value</u>
Weight (lb.)	141
Cost (\$10 ⁴)	\$233
Power (watts)	74

5.6 Analysis and Discussion of Results

The decision problem presented in Section 5.5 was originally solved via a "multirisk" programming model [6]. This analysis concept is based on the determination of the alternative decision solutions which minimize the probabilities that the decision maker's objectives and constraints will not be satisfied. The "best" subset of m decision alternatives from a possible set of n candidates is selected such that the problem's objectives and constraints are satisfied with minimum risk.

The multirisk programming model is designed to solve multiple criteria decision problems, assuming that decision makers strive to achieve or satisfy goals rather than attempting to optimize them. The decision maker may incorporate the concept of "fuzziness" [66] into the analysis by specifying a range of deviation allowable for the goals and constraints of the problem.

The multirisk programming model is a stochastic analysis technique for solving multiple criteria decision problems. Problem formulation may include both rigid and stochastic goals and constraints. The rigid goals and constraints have deterministic coefficients, while the stochastic goals and constraints have stochastic coefficients. The model assumes that the range of deviation is a random variable, while the stochastic parameters of the objectives and constraints are normally distributed independent random variables. The "best" solution to the multirisk programming model is the program which minimizes the risk of not achieving the goals and constraints of the problem, or maximizes the probability that the goals and constraints are achieved.

In the multirisk programming analysis of the project selection example (Section 5.5), uncertainties were incorporated for each coefficient. The attribute data for the proposed candidates is shown in Table V.13.

The multirisk programming model employs an enumerative search to identify the "best" solution. Table V.14 displays the results of the multirisk programming model and the fuzzy linear integer programming/exchange heuristic model. Both fuzzy linear integer programming/exchange heuristic model and the multirisk programming model of multiple criteria decision making provide effective mathematical modeling techniques for the analysis of management decisions which are fuzzy in nature.

The problem formulation of the fuzzy linear programming/exchange heuristic model was presented in Section 5.2. This multiple objective analysis provides the decision maker "leeway" in modeling phenomena of a vague or ill-defined nature. This "leeway" in the model is a result of utilizing fuzzy sets to describe those objective functions and constraints that are imprecisely defined. Through the use of a fuzzy set operator, the "min" operator, an optimal decision in the fuzzy environment is obtained. This optimal decision is defined as the point which maximizes the membership function of the fuzzy set formed through the intersection of those fuzzy sets representing the various objective functions and constraints. The exchange heuristic in the model seeks to obtain the best attainable integer solution given that the optimal point defined above is not integer valued.

Table V.13. Attribute Data for Proposed Candidates

<u>Attribute</u>	<u>Subsystem A</u>			<u>Subsystem B</u>			<u>Subsystem C</u>			
	<u>1</u>	<u>2</u>	<u>3</u>	<u>4</u>	<u>5</u>	<u>6</u>	<u>7</u>	<u>8</u>	<u>9</u>	<u>10</u>
Weight (lb.)	32	57	19	95	107	61	48	23	10	15
Cost (\$10 ⁴)	120	95	160	64	67	96	119	42	36	70
Reliability	0.97	0.94	0.99	0.89	0.90	0.94	0.96	0.98	0.97	0.99
Power (watts)	21	35	10	60	83	27	50	12	7	16

Uncertainties in weights = $\pm 10\%$.

Uncertainties in costs = $\pm 15\%$.

Reliabilities are assumed constant.

Uncertainties in power requirements = $\pm 5\%$.

Table V.14. Comparison of Results

<u>System Attribute</u>	<u>Fuzzy LP/Heuristic Model</u>	<u>Multirisk Programming Model</u>	
		<u>Expected Value</u>	<u>Standard Deviation</u>
Weight	141 lbs.	141 lbs.	4.3 lbs.
Cost	\$2,330,000	\$2,330,000	\$106,000
Reliability	0.866	0.866	0.0
Power	74 watts	74 watts	1.1 watts

One practical advantage of modeling with either of these techniques is that the decision situation does not have to be defined in a precise manner. In the fuzzy programming approach, the decision maker specifies ranges of acceptability for those objective functions and constraints represented by fuzzy sets. In the multirisk programming model, the decision maker specifies ranges of deviation. In both models, the decision maker is given greater flexibility than would be available in a classical mathematics approach.

A major advantage of fuzzy linear programming is the ease with which it can be formulated and solved on numerous mathematical programming systems. The exchange heuristic and the multirisk programming techniques both require utilizing a specific computer program which may not be readily available.

The major advantage of the multirisk programming model is its ability to analyze multiple criteria decision problems which are characterized by nondeterministic coefficients for the various objective functions. This model provides the decision maker leeway in defining his aspiration levels, as well as in stating precisely the terms of the objective functions.

The enumerative search technique employed in the multirisk programming model is impractical for large scale problems both in computer storage requirements and necessary CPU time [6]. The fuzzy linear programming model utilizes whatever mathematical programming system that is available to the user to solve mixed integer linear programming problems. The computer storage requirement is, therefore, programming package dependent, as is the CPU time necessitated.

5.7 Computer Program

A computer program implementing the fuzzy linear integer programming and the exchange heuristic phases of the algorithm was developed in FORTRAN IV for the IBM 370/3033 computer system. Standard FORTRAN language was employed to permit relative ease of adaptation of the computer model for use on other computer systems. The amount of internal storage necessitated on the IBM 370/3033 was 280,000 bytes. The amount of storage necessary is due to the requirements of the MPSX system. The computer program is currently dimensioned for the comparison of twenty-five alternatives. This program size could be enlarged by redimensioning the program not to exceed the MPSX variable limit. The exchange heuristic program is capable of solving problems of size one hundred and fifty constraints with one hundred and fifty decision variables with 280,000 bytes of storage. The CPU time to execute Example 1 was 2 seconds, while 3 seconds of CPU time was required for Example 2.

5.8 Computer Code Description

The fuzzy linear integer programming computer code and the exchange heuristic computer code are listed in Appendix B along with the definition of all input data. The computer codes will be described in three sections:

- i. Fuzzy Linear Programming Transformation program.
- ii. IBM MPSX/Mixed Integer Programming Control program.
- iii. Exchange Heuristic program.

The fuzzy linear integer programming transformation program is composed of a single main program. This program reads the input data

and transforms the objective functions into their fuzzy transformation as appropriate. Temporary data sets are created to be used as data input for the MPSX/Mixed Integer Programming Optimization Technique or the exchange heuristic program as appropriate.

The MPSX/Mixed Integer Programming Control program is an advanced usage example [63] of the IBM MPSX/MIP technique. This control program optimizes the continuous problem, then solves the mixed integer problem. A more in-depth discussion of the MPSX mathematical programming system is presented in Appendix A.

The exchange heuristic code consists of a main program and seven subroutines. The code is a modified version of Petersen's [42] heuristic algorithm, with subsequent modifications by Bouillot and Smith [64, 65]. The description of this code is as follows:

Main Program. The main program reads the input data from the temporary data set created in the fuzzy linear programming transformation program. This program calculates the R_j values and the ratio T_j/R_j for each variable. It maintains the list of selected and nonselected variables and determines those exchanges to be executed. It calls the various subroutines in the proper sequence required to conduct the exchange operations. It formats and writes all output data as required.

Subroutine Rank. This subroutine initially ranks the variables in both the sets of selected and nonselected variables.

Subroutine Impvmt. This subroutine maintains the best solution achieved that is both advantageous and feasible.

Subroutine Feasbl. This subroutine examines the feasibility of the exchange under consideration.

Subroutine Exchge. This subroutine executes the exchange operations.

Subroutine Achvmt. This subroutine calculates the gain in the membership function as a result of an exchange.

Subroutine OBJ. This subroutine evaluates the various fuzzy objective functions to determine the lambda value.

Subroutine FTOBJ. This subroutine calculates the fitback solution.

CHAPTER 6

SUMMARY, CONCLUSIONS, AND SUGGESTIONS FOR FURTHER RESEARCH

In this chapter, the work presented in this thesis is summarized, and a few conclusions are drawn about the fuzzy capital budgeting model and the general applicability of fuzzy programming.

6.1 Summary and Conclusions

The model of the capital budgeting problem explored in this work is a combined application of the models developed by Zimmermann [11] and Petersen [42]. The fuzzy linear integer programming approach to management decisions is designed to study decision problems involving multiple goals and constraints, some of which are of a vague or ill-defined nature. The method is founded on the theory of fuzzy sets. Fuzzy sets are utilized to model phenomena of an ill-defined nature which cannot be described adequately in classical mathematical terms. The analysis seeks to permit the human mind to utilize its capabilities to the fullest extent, while utilizing the computational efficiency of the computer to perform those operations which the human mind cannot adequately accomplish. This analysis allows the individual decision maker to make small judgmental decisions (what the human mind does best), while allowing the computer to solve large linear programming problems incorporating these judgmental decisions. Since the solution to the fuzzy linear integer programming problem may not be in the form $(0,1)$, a modified form of Petersen's exchange heuristic for the capital

budgeting problem was employed. This exchange heuristic seeks the "best" attainable solution, given that no integer solution exists.

The model is designed to solve the general problem in which the "best" subset of m alternatives is selected from a candidate set of n possible decision alternatives, such that the membership function of the fuzzy set of the decision is maximized. The model provides a great deal of flexibility to the user in formulating problems for analysis. The availability of solution algorithms and computer solution systems which are readily compatible with the fuzzy linear integer programming problem formulation allows the user to realize a computational solution with ease.

This analysis has applicability for a broad range of decision problems involving the selection of entities from among numerous alternative possibilities, such as equipment purchases, route selection, or investment selection. The example presented in Chapter 5 successfully analyzed the selection of alternative subsystems in achieving system design requirements, while satisfying stated cost restrictions.

6.2 Suggestions for Further Research

The work described in this thesis can be extended in several different directions. The integer solution technique utilized in this work was the MPSX/Mixed Integer Programming System. Although this mathematical programming system readily yields a solution to the problem, the availability of MPSX is not universal. The development and use of an integer programming computer code in standard FORTRAN would greatly enhance the ease with which the model could be adapted to other computers.

The popularity of multiple objective analysis via goal programming could be the catalyst of another extension. The formulation of the fuzzy objectives and constraints into a goal programming analysis would be an extremely interesting development. Goal programming is an extremely robust optimization technique, which is viewed as a practical and natural representation of a wide variety of real-world problems. In combining the fuzzy programming approach of optimizing humanistic systems which are by nature vague and ill-defined, with the practicality of goal programming, an optimization methodology may result which presents a realistic perspective of management decision making.

Field experimentation with fuzzy programming models of management decision making would be desirable. Zimmermann [58] has conducted numerous experiments to analyze the viability of modeling decision makers via the concepts inherent in fuzzy programming. Applications of fuzzy programming include personnel management and determination of credit worthiness in the banking industry [58], media selection [11], and the sizing of a truck fleet [9].

The fuzzy set operator used in the fuzzy linear integer programming model is the "min" operator. Zimmermann and Hamacher [58] have experimented on the applicability of other operators in the optimization of management decisions. These operators include the product operator, the algebraic sum operator, the max operator, both the arithmetic and geometric mean operators, and the gamma operator.

Certainly, many other areas for further research exist. However, these few are listed to provide the reader some idea as to where additional research might begin.

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As in most integer programming techniques, long computational time is necessary to determine the optimal integer solution.

The IBM program descriptions [63] provide an in-depth discussion of the MPSX system. In addition, numerous sample problems are presented. The examples are detailed from data input through sample outputs. The Mixed Integer Programming program description is especially instructive. It provides sample MPSX control programs and presents an excellent discussion of the mixed integer programming procedure.

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APPENDIX A

A BRIEF LOOK AT MPSX

In this appendix, the IBM Mathematical Programming System Extended (MPSX), linear programming and mixed integer programming capabilities are summarized.

A.1 MPSX System

Mathematical Programming System Extended (MPSX) is composed of a set of procedures all operating under the direction of a user specified MPSX control program. Through the MPSX control program, the user specifies the sequence of steps to be executed in solving a mathematical programming problem.

The user is able to augment MPSX with procedures written in the FORTRAN language through the use of the Read Communications Format (READCOMM) feature of MPSX. Through the use of FORTRAN CALL statements, the READCOMM subroutine is accessed. This subroutine acts as an interface between the MPSX control program and the FORTRAN procedures.

The user is capable of executing all of the MPSX capabilities through the use of the MPSX control programs and the READCOMM procedures.

A.2 Linear Programming Procedure

The MPSX strategy for solving a linear programming problem is the ordered execution of a series of the MPSX procedures. The user specifies the solution strategy to MPSX, via the MPSX control language.

The linear programming procedures of MPSX use the bounded variable/product form of the inverse/revised simplex. The simplex method is based upon the fact that if there are m constraints which are linearly independent, then there is a set of m columns (variables) which are also linearly independent. The right-hand side values can be expressed in terms of the m columns called a basis. The simplex

method employs these basic solutions by exchanging one column from the basis with one column not in the basis on each iteration, until a solution is realized that satisfies the feasibility criteria. This solution is termed a basic feasible solution.

The simplex method proceeds by examining the basic feasible solutions, to find one that satisfies the requirement that the objective function value be maximized or minimized.

A.3 Mixed Integer Programming Procedure

The Mixed Integer Programming capability of MPSX is an extension of the linear programming procedure of MPSX. It provides the user the capability to solve linear programming problems composed of both integer and continuous variables. This analysis is appropriate for the fuzzy linear programming approach to the capital budgeting problem. Since the solution must be of the form $(0,1)$, the variables representing the investment possibilities must be integer values, while the value of the membership function is a continuous variable.

The MPSX mixed integer linear programming problem is performed in two stages. First, the problem is solved considering all integer variables as being continuous. The problem is solved by the linear programming capability of MPSX. The solution to this problem is termed the optimal continuous solution.

The second stage is to solve the problem for the optimal integer solution. The search for an integer solution starts from the optimal continuous solution and proceeds using the branch and bound technique. The search continues until the optimal integer solution is determined.

APPENDIX B

Fuzzy Linear Integer Programming Via
MPSX/MIP Computer Code, Variable
Definition Guide, User's Guide,
and Sample Output

HFADLY M .. HASP-II***.....START JOB 1718.....RF229841.....HEADLY M
 HFADLY M .. HASP-II***.....START JOB 1718.....RF229841.....HFADLY M

TIME: 15:33:21 DATE: 03/17/80 H A S P S Y S T E M L O G

```

**NP229841 JOB '04771,T=0050,K=01500,S=280,          ' , 'HEADLY M          ' , ' AA
***FULLSKIPS
// EXEC SETTAF,TTAIN=TH,FORMS=16
**INCLUDE MGH01.$TODAY
IEF142I - STEP WAS EXECUTED - COND CODE 0000
IEF373I STEP /SETTAF / START 80077.1533
IEF374I STEP /SETTAF / STOP 80077.1533 CPU          0MIN 00.03SEC MAIN          8K LCS          OK
// EXEC PGM
//SYSIN DD *
//DATA.FT45F001 DD UNIT=SYSDA,DSN=88NAMFS,SPACE=(3200,(50,5)),
// DISP=(NEW,PASS),DCB=(RECFM=FB,LRECL=80,BLKSIZE=3200)
//DATA.INPUT DD *
IEF142I - STEP WAS EXECUTED - COND CODE 0000
IEF373I STEP /DATA / START 80077.1533
IEF374I STEP /DATA / STOP 80077.1533 CPU          0MIN 00.37SEC MAIN          280K LCS          OK
// EXEC LP
//SOURCE.INPUT DD *
**INCLUDE MGH01.$MISSY
IEF142I - STEP WAS EXECUTED - COND CODE 0000
IEF373I STEP /SOURCE / START 80077.1533
IEF374I STEP /SOURCE / STOP 80077.1534 CPU          0MIN 00.70SEC MAIN          280K LCS          OK
//DATA.INPUT DD UNIT=SYSDA,DSN=88NAMFS,DISP=(OLD,DELETE)
IEF142I - STEP WAS EXECUTED - COND CODE 0000
IEF373I STEP /DATA / START 80077.1534
IEF374I STEP /DATA / STOP 80077.1534 CPU          0MIN 00.47SEC MAIN          280K LCS          OK
IEF375I JOB /RF229841/ START 80077.1533
IEF376I JOB /RF229841/ STOP 80077.1534 CPU          0MIN 01.57SEC

```

 FUZZY INTEGFF (0-1) LINEAR PROGRAMMING

SOLUTION TECHNIQUE FOR THE MULTIOBJECTIVE

CAPITAL BUDGETING PROBLFM VIA MPSX/MIP

PREPARED BY CPT. MICHAEL G. HEADLY

 ADDITIONAL INFORMATION CONCERNING FUZZY PROGRAMMING

MAY BE OBTAINED IN THE FOLLOWING REFERENCES

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```

** ** X6(I,J): COEFFICIENT OF THE JTH VARIABLE IN THE
** ** ITH TYPE3 FUZZY OBJECTIVE FUNCTION.
**
** ASP3(I): ASPIRATION LEVEL ITH TYPE3 FUZZY OBJECTIVE
**
** STINT(I): LOWEST ADMISSABLE VALUE OF THE ITH
** TYPE3 FUZZY CEJECTIVE.
**
** UTINT(I): GREATEST ADMISSABLE VALUE OF THE ITH
** TYPE3 FUZZY CEJECTIVE.
**
** INTERNAL VARIABLES
** *****
**
** ROW(I,J): CCEFFICENT OF THE JTH VARIABLE IN THE ITH
** CONSTRAINT/OBJECTIVE FORMULATED FOR INPUT INTO MPSX.
**
** RHS(I): RIGHT-HAND SIDE VALUE OF THE ITH CONSTRAINT/
** OBJECTIVE FORMULATED FOR INPUT INTO MPSX.
**
** MPSX INPUT COMMANDS
** *****
**
** INTORG AND INTEND: MPSX COMMANDS WHICH IDENTIFY THOSE
** VARIABLES WHICH MUST BE INTEGER VALUED.
**
** UPBOUND: MPSX COMMAND WHICH ESTABLISHES THE
** UPPER BOUND OF THOSE VARIABLES THAT ARE BOUNDED.
**
** *****

```



```
C
C      ***
C      ** IF(KTYPE2.EQ.0) GO TO 31
C      DO 51 JA=1,KTYPE2
C      READ 2,(X2(IA,I),I=1,NVAR)
C      READ 2,(RHS2(I),I=1,KTYPE2)
C      DO 300 I=1,KTYPE2
C      M=M+1
C      DO 400 J=1,NVAR
C      ROW(M,J)=X2(I,J)
C      RHS(M)=RHS2(I)
C      CONTINUE
C      CONTINUE
C      ***
C      *** RIGID CONSTRAINT TYPE 3 (=
C      ***
C      ***
C      *** M=KTYPE1+KTYPE2
C      31 IF(KTYPE3.EQ.0) GO TO 32
C      DO 52 IB=1,KTYPE3
C      READ 2,(X3(IB,I),I=1,NVAR)
C      READ 2,(RHS3(I),I=1,KTYPE3)
C      DO 500 I=1,KTYPE3
C      M=M+1
C      DO 600 J=1,NVAR
C      ROW(M,J)=X3(I,J)
C      RHS(M)=RHS3(I)
C      CONTINUE
C      CONTINUE
C      ***
C      ***
C      *** FUZZY CONSTRAINT TYPE 1
C      ***
C      ***
```

ARMY MILITARY PERSONNEL CENTER ALEXANDRIA VA F/G 5/3
AN ANALYSIS OF THE MULTIPLE OBJECTIVE CAPITAL BUDGETING PROBLEM--ETC(U)
MAY 80 M G HEADLY

F/G 5/3

NL

2 of 2

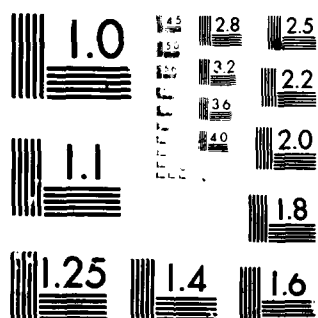
ΔM₁
ΔM₂ 1.75, 1.75

END

DATE _____

6-8

DTIC



MICROCOPY RESOLUTION TEST CHART
NATIONAL BUREAU OF STANDARDS-1963-A

```

32  M=K*TYPE1+K*TYPE2+K*TYPE3
   IF (I*TYPE1.EQ.0) GC TO 33
   DO 60 ID=1, I*TYPE1
   READ 2, (X4(ID,I), I=1, NVAR)
   READ 2, (ASP1(I), TINT1(I), I=1, I*TYPE1)
   DO 700 I=1, I*TYPE1
   M=M+1
   DO 800 J=1, NVAR
   POW(M,J)=-X4(I,J)/TINT1(I)
   EHS(M)=(TINT1(I)-ASP1(I))/TINT1(I)
800 CONTINUE
700 CONTINUE
C ***
C ***
C *** FUZZY CONSTRAINT TYPE 2
C ***
C ***
33  M=K*TYPE1+K*TYPE2+K*TYPE3+I*TYPE1
   IF (I*TYPE2.EQ.0) GC TO 34
   DO 53 IF=1, I*TYPE2
   READ 2, (X5(IF,I), I=1, NVAR)
   READ 2, (ASP2(I), TINT2(I), I=1, I*TYPE2)
   DO 900 I=1, I*TYPE2
   M=M+1
   DO 1000 J=1, NVAR
   POW(M,J)=X5(I,J)/TINT2(I)
   EHS(M)=(TINT2(I)+ASP2(I))/TINT2(I)
1000 CONTINUE
900 CONTINUE
C ***
C ***
C *** FUZZY CONSTRAINT TYPE 3
C ***
C ***
34  M=K*TYPE1+K*TYPE2+K*TYPE3+I*TYPE1+I*TYPE2

```

```

IF (ITYPE3.EQ.0) GO TO 35
DO 54 IG=1, ITYPE3
  READ 2, (X6(IG,I), I=1, NVAR)
  READ 2, (ASP3(I), STINT(I), UTINT(I), I=1, ITYPE3)
DO 1100 I=1, ITYPE3
  M=M+1
MCOUNT=M
DO 1200 J=1, NVAR
  FOR (A,J)=-X6(I,J)/STINT(I)
  RHS(M)=- (ASP3(I)-STINT(J))/STINT(I)
  M=M+1
  ROW (M,J)=X6(I,J)/UTINT(I)
  RHS(M)= (UTINT(I)+ASP3(I))/UTINT(I)
  IF (J.EQ.NVAR) GO TO 1200
  M=M-1
1200 CONTINUE
  KKJ=MCOUNT
  KJJ=KKJ+1
  KKJ=0
1100 CONTINUE
  ***
  *** THE REMAINDER OF THE PROGRAM WRITES
  ***
  *** THE TRANSFORMED FUNCTIONS INTO
  ***
  *** THE TEMPORARY DATA SET FOR INPUT INTO
  *** THE MPSX/MIP ROUTINE.
  ***
35 WRITE(45,8)
8  FORMAT('NAME',10X,'MAX')
9  WRITE(45,9)
9  FORMAT('ROWS')
10 WRITE(45,10)
10 FORMAT(2X,'N',1X,'OBJ')
  M=0

```

```

11 IF (KTYPE1.EQ.0) GO TO 36
DO 2000 I=1,KTYPE1
M=I
M=M+100000
WRITE(45,11) M
FORMAT(2X,'L',1X,16,T5,'ROW')
M=M-100000
2000 CONTINUE
36 M=KTYPE1
IF (KTYPE2.EQ.0) GO TO 37
DO 3000 I=1,KTYPE2
M=M+1
M=M+100000
WRITE(45,12) M
FORMAT(2X,'G',1X,16,T5,'ROW')
M=M-100000
3000 CONTINUE
37 M=KTYPE1+KTYPE2
IF (KTYPE3.EQ.0) GO TO 38
DO 4000 I=1,KTYPE3
M=M+1
M=M+100000
WRITE(45,13) M
FORMAT(2X,'E',1X,16,T5,'ROW')
M=M-100000
4000 CONTINUE
38 M=KTYPE1+KTYPE2+KTYPE3
IF (KTYPE1.EQ.0) GO TO 39
DO 5000 I=1,KTYPE1
M=M+1
M=M+100000
WRITE(45,11) M
M=M-100000
5000 CONTINUE
39 M=KTYPE1+KTYPE2+KTYPE3+KTYPE1

```

```

IF (IYPE2.EQ.0) GO TO 40
DO 6000 I=1,IYPE2
M=M+1
M=M+100000
WRITE(45,11) M
M=M-100000
CONTINUE
40 N=KTYPE1+KTYPE2+KTYPE3+IYPE1+IYPE2
IF (IYPE3.EQ.0) GO TO 41
K=M+IYPE3*2
M=M+1
M=M+100000
WRITE(45,11) M
M=M-100000
IF (M.NE.K) GO TO 42
41 WRITE(45,14)
14 FORMAT('COLUMNS')
15 WRITE(45,15)
15 FORMAT(4X,'LAMBDA',4X,'CEJ',9X,'1.0')
L=IYPE1+IYPE2+IYPE3*2
M=KTYPE1+KTYPE2+KTYPE3
DO 8000 I=1,L
M=M+1
M=M+100000
WRITE(45,16) M
16 FORMAT(4X,'LAMBDA',4X,I6,T15,'RCM',7X,'1.0')
M=M-100000
CONTINUE
45 WRITE(45,45)
45 FORMAT(4X,'INIOR3',4X,'...MARKER...',17X,'...INTORG...')
H=0
L=0
L=KTYPE1+KTYPE2+KTYPE3+IYPE1+IYPE2+IYPE3*2

```



```

DO 8050 J=1,NVAR
DO 5000 I=1,L
M=M+1
IF (M.FO.L+1) M=1
N=J
K=M
N=N+1000
K=K+100000
WRITE(45,17) N,K,ROW(M,J)
FORMAT(4X,I4,T5,'X',9X,I6,T15,'POW',7X,F8.4)
N=N-1000
K=K-100000
CONTINUE
CONTINUE
WRITE(45,46)
FORMAT(4X,'INTEND',4X,'...MARKER...',17X,'...INTEND...')
WRITE(45,18)
FORMAT('RHS')
L=KTYPE1+KTYPE2+KTYPE3+ITYPE1+ITYPE2+ITYPE3*2
DO 9050 I=1,L
M=I
N=N+100000
WRITE(45,19) N,RHS(I)
FORMAT(4X,'RHS',7X,I6,T15,'POW',7X,F8.4)
N=N-100000
CONTINUE
WRITE(45,47)
FORMAT('BOUNDS')
WRITE(45,49)
FORMAT(1X,'UP',1X,'UPBCUND',3X,'LAMBDA',4X,'1.0')
DO 9075 J=1,NVAR
N=J
N=N+1000
WRITE(45,48) N
FORMAT(1X,'UP',1X,'UPBCUND',3X,I4,T15,'X',10X,'1.0')
N=N-1000

```

```

9075 CONTINUE
      WRITE(45,20)
20  FORMAT('FNDATA')
      1  FORMAT(7I5)
      2  FORMAT(9F8.4)
      STOP
      END

```

ECHOED= 1775

OBJECT CODE= 15048 BYTES,ARRAY AREA= 18600 BYTES,TOTAL AREA AVAILABLE= 178176

NUMBER OF ERRORS= 0, NUMBER OF WARNINGS= 0, NUMBER OF EXTENSIONS=

= 0.22 SEC,EXECUTION TIME= 0.08 SEC,

..MPSX-VIM7.. CONTROL PROGRAM COMPILE. MPSX RELEASE 1 MOD LEVEL 7

0001	PROGRAM	
0002	INITIALZ	
0096	MOVE(XDATA,'MAX')	
0097	MOVE(XPDNAME,'EUDGET')	
0098	CONVERT	
0099	SETUP('BOUND','UPBOUND','MAX')	
0100	MOVE(XOBJ,'CBJ')	
0101	MOVE(XRHS,'RHS')	
0102	OPTIMIZE	
0411	SOLUTION	
0412	SAVE('NAME','CETC')	
0413	INIMIX	
0451	MIXSTART('CCST')	
0452	XMXDROP=0.0	
0453	CT=0	
0454	MVADR(XDOPRINT,INT)	
0455	MIXFLOP	
0456	MIXSAVE('NAME','TRER')	STOP
0457	MIXSTATS('NODES')	
0458	FXIT	
0459	SOLUTION	INT
0460	XMXDROP=0.0	
0461	CT=CT+1	
0462	IF(CT.EQ.5,STOP)	
0463	CONTINUE	
0464	LC(0)	CT
0465	PEND	

..MPSX-V1M7.. EXECUTOR. MPSX RELEASE 1 MCD LEVEL 7

CONVERT MAX TO BUDGET

TIME = 0.00

1- F0ES SECTION.

0 MINOR ERROR(S) - 0 MAJOR ERROR(S).

2- COLUMNS SECTION.

0 MINOR ERROR(S) - 0 MAJOR ERROR(S).

3- RHS'S SECTION.

RHS

0 MINOR ERROR(S) - 0 MAJOR ERROR(S).

5- BOUNDS SECTION.

UBOUND

0 MINOR ERROR(S) - 0 MAJOR ERROR(S).

PROBLEM STATISTICS

12 LP ROWS, 23 VARIABLES, 92 LP ELEMENTS, DENSITY = 33.33

10 INTEGER VARIABLES

THESE STATISTICS CONTAIN ONE SLACK VARIABLE FOR EACH ROW

0 MINOR ERROR(S), 0 MAJOR ERROR(S).

..MPSX-V147.. EXECUTED. MPSX RELEASE 1 RCD LVFL 7

SFTUP BUDGET

TIME = 0.00

BOUND = UNBOUND
 MAX
 SCALF

MATRIX1 ASSIGNED TO MATRIX1
 MATRIX2 ASSIGNED TO MATRIX2

ETA1 ASSIGNED TO ETA1

SCRATCH1 ASSIGNED TO SCRATCH1
 SCRATCH2 ASSIGNED TO SCRATCH2

WORK ASSIGNED TO WORK

MAXIMUM PRICING NOT REQUIRED - MAXIMUM POSSIBLE 7

NO CYCLING

POOLS	NUMBER	SIZE	CCNE
P. REG-BITS MAP			72
INTEGER TABLES			1568
BOUND VECTOR			120
WORK REGIONS	9	120	1080
MATRIX BUFFERS	3	600	1800
ETA BUFFERS	3	2304	6912

POOLS (LOG.VAR.)	TOTAL	NORMAL	FREE	FIXED	BOUNDED
COLUMNS (STR.VAR.)	12	6	1	5	0
	11	0	0	0	11

10 COLUMNS ARE INTEGER
 92 ELEMENTS - DENSITY = 33.33 - 3 MATRIX RECORDS (WITHOUT RHS'S)

WRITE

TIME = 0.00

OPTIMIZE SYSTEM MACRO CALLED

CFASH TIME 0.00 MINS.
 INFASIBILITIES 8 AT START
 8 AFTER PASS A

INVERT CALLED	TIME	0.00	CURRENT INVERSE	----	FTA-VICTORS1	ELEMENTS1
PASS	----	NO. OF ROWS12	LOGICALS12	STRUCTURALS0	ELEMENTS
INVERSE	--	NUCLEUS0	TRANSPICED0	FTA-VICTORS1	ELEMENTS

0 AFTER PASS B
 FEASIBLE

NEGDI

TIME = 0.00 MINS.

PRIMAL OBJ = OBJ PHS = RHS

TIME = 0.00 MINS.
 SCALE = 1.00000- PRICING 7

```

ITER  NUMBER  VECTOR  VFCIOR  REDUCED  FUNCTION
NUMBER  NOOPT   OUT    IN    COST    VALUE
4      10      1      11      13      1.00000-  .13340
4      11      1      16      14      .33182-  .13340
INVERT DEMANDED AFTER 10 MAJOR/ 10 MINOR ITERATIONS - CLOCK CONTROL

INVERT CALLED      TIME 0.00  CURRENT INVERSE ---- ETA-VECTORS .....7  ELEMENTS .....44
BASIS ---- NO.OF ROWS .....12  LOGICALS .....7  STRUCTURALS .....5  ELEMENTS .....42
INVERSE -- NUCLEUS .....0  TRANSFORMED .....0  ETA-VECTORS .....6  ELEMENTS .....36

```

NEGDJ

TIME = 0.00 MINS.

PRIMAL OBJ = OBJ RHS = RHS

TIME = 0.00 MINS. PRICING 7
SCALE = 1.00000-

```

ITER  NUMBER  VECTOR  VFCIOR  REDUCED  FUNCTION
NUMBER  NOOPT   OUT    IN    COST    VALUE
4      12      1      2      17      .46110-  .50945
4      13      1      6      18      .40714-  .50945
4      14      1      22     23      2.50491-  .50945
4      15      1      8      15      .61259  .58071
4      16      1      20     22      .54614-  .58071
INVERT DEMANDED AFTER 5 MAJOR/ 5 MINOR ITERATIONS - CLOCK CONTROL

INVERT CALLED      TIME 0.00  CURRENT INVERSE ---- ETA-VECTORS .....11  ELEMENTS .....85
BASIS ---- NO.OF ROWS .....12  LOGICALS .....4  STRUCTURALS .....8  ELEMENTS .....62
INVERSE -- NUCLEUS .....6  TRANSFORMED .....2  ETA-VECTORS .....14  ELEMENTS .....65

```

NEGDJ

TIME = 0.00 MINS.

PRIMAL OBJ = OBJ RHS = RHS

..MPX-V17.. EXECUTOR. MPX RELEASE 1 MOD LEVEL 7

TIME = 0.00 MINS. PENDING 7

SCALE = 1.00000-

OPTIMAL SOLUTION

..MPX-V17.. EXECUTOR. MPX RELEASE 1 MOD LEVEL 7

SOLUTION (OPTIMAL)

TIME = 0.00 MINS. ITERATION NUMBER = 16

...NAME... ...ACTIVITY... DEFINED AS

FUNCTIONAL .58071 OBJ

RESTRAINTS FHS

BOUNDS... UFGUND

17.. EXECUTOR. MPX RELEASE 1 MOD LEVEL 7

1 - ROWS

...ROW...	AT	...ACTIVITY...	SLACK ACTIVITY	..LOWER LIMIT.	..UPPER LIMIT.	DUAL ACTIVITY
GRJ	ES	.58071	.58071-	NONE	NONE	1.00000
ROW001	LL	.16250-	.	.16250-	NONE	9.44695
ROW002	EQ	1.00000	.	1.00000	1.00000	2.03095
ROW003	EQ	1.00000	.	1.00000	1.00000	1.96510
ROW004	EQ	1.00000	.	1.00000	1.00000	.80932
ROW005	EQ	1.00000	.	1.00000	1.00000	.09136-
ROW006	EQ29727-
ROW007	UL	4.00000	.	NONE	4.00000	.95977-
ROW008	ES	4.53436-	.53436	NONE	4.00000-	.
ROW009	BS	10.61078	.18922	NONE	11.00000	.
ROW010	UL	2.33330-	.	NONE	2.33330-	.04023-
ROW011	ES	9.32265	1.67735	NONE	11.00000	.

..MPSX-V117.. EXECUTOR. MPSX RELEASE 1 MOD LPVFL 7

SECTION 2 - COLUMNS

NUMBER	COLUMN.	AT	...ACTIVITY...	..INPUT COST..	..LOWER LIMIT.	..UPPER LIMIT.
13	LAMBDA	BS	.58071	1.00000	.	1.00000
14	X001	BS	.06172	.	.	1.00000
15	X002	BS	.93828	.	.	1.00000
16	X003	LL	.	.	.	1.00000
17	X004	BS	.32810	.	.	1.00000
18	X005	BS	.06172	.	.	1.00000
19	X006	BS	.61018	.	.	1.00000
20	X007	LL	.	.	.	1.00000
21	X008	HL	1.00000	.	.	1.00000
22	X009	BS	.	.	.	1.00000
23	X010	BS	.	.	.	1.00000

..MPSX-V117.. EXECUTOR. MPSX RELEASE 1 MOD LPVFL 7

MYSTART - TIME = 0.00

COST

MYSTOP - TIME = 0.00

I NODES DISPOSED IF ESTIMATION BEYOND -100000R+76-
 I NODES DROPPED IF FUNCTIONAL BEYOND .
 I REACH NODE 1 - FUNCTIONAL = .5807 ESTIMATION = .5807 ITER 10
 I FIRST BRANCH - VARIABLE 17 = .3281 FORCED TOWARDS UPPER ACTIVITIES

ITER	NUMBER	VECTOR	VECTOR	REDUCED	FUNCTION	PARAM
NUMBER	NONOPT	OUT	IN	COST	VALUE	VALUE
M	17	0	17	.21563	.58071	.32810
M	18	0	22	.13532-	.58071	.32810
M	19	0	18	.13981	.56528	.36271
M	20	0	19	.09256	.37256	.58091

BRANCH DROPPED

I RESTORE NODE 1 - FUNCTIONAL = .5807 ESTIMATION = .5807
 I SECOND BRANCH - VARIABLE 17 = .3281 FORCED TOWARDS LOWER ACTIVITIES

M	21	0	17	.01927-	.58071	.67190
M	22	0	22	.03390	.50633	.92741

I REACH NODE 2 - FUNCTIONAL = .5506 ESTIMATION = .55058 ITER 22
 I BEST POSSIBLE SOLUTION IS NOT .55058
 I FIRST BRANCH - VARIABLE 14 = .6967 FORCED TOWARDS LOWER ACTIVITIES

M	23	0	14	.24577	.55058	.30328
M	24	0	11	.50121	.53262	.59577

I REACH NODE 3 - FUNCTIONAL = .3190 ESTIMATION = .31897 ITER 24
 I RESTORE NODE 2 - FUNCTIONAL = .5506 ESTIMATION = .55058
 I SECOND BRANCH - VARIABLE 14 = .6967 FORCED TOWARDS UPPER ACTIVITIES

M	25	0	14	.05501-	.55058	.69672
M	26	0	12	2.22737	.54413	.82795

BRANCH DROPPED

```

1 BEST POSSIBLE SOLUTION IS NOW      .31897
1 RESTORE NODE      3 - FUNCTIONAL = .3190 ESTIMATION = .31897
1 FIRST BRANCH - VARIABLE 15 = .6591 FORCED TOWARDS LOWER ACTIVITIES
4 27 0 15 11 .39711 .31897 .34091
BRANCH DROPPED

```

```

1 RESTORE NODE      3 - FUNCTIONAL = .3190 ESTIMATION = .31897
1 SECOND BRANCH - VARIABLE 15 = .6591 FORCED TOWARDS UPPER ACTIVITIES
4 28 0 15 22 .00205- .31897 .65909
4 29 0 21 20 .09723 .31699 .82233
4 30 0 22 21 .03283- .31699 .82233

```

```

1 REACH NODE      4 - FUNCTIONAL = .2841 ESTIMATION = .28406
1 BEST POSSIBLE SOLUTION IS NOW      .28406
1 FIRST BRANCH - VARIABLE 19 = .8326 FORCED TOWARDS LOWER ACTIVITIES

```

```

INVERT CALLED      TIME 0.00 CURRENT INVERSE ---- ETA-VFACTORS ....21 ELEMENTS ...138 RECC
BASIS ---- NO.OF ROWS ....12 LOGICALS .....5 STRUCTURALS .....7 ELEMENTS ....56
INVERSE -- NUCLEUS .....5 TRANSFORMED .....1 ETA-VECTORS .....13 ELEMENTS ....58 RECC

```

MIXFLOW - TIME = 0.00

```

ITER NUMBER VECTOR VECTOR VECTOR VECTOR VECTOR VECTOR VECTOR VECTOR
NUMBER NOOPT CUT 19 9 11 11 11 11 11 11 11 11 11 11 11 11 11 11 11 11
31 0 19 9 11 11 11 11 11 11 11 11 11 11 11 11 11 11 11
32 0 19 9 11 11 11 11 11 11 11 11 11 11 11 11 11 11 11
BRANCH DROPPED

```

```

-----
I RESTOPED MODE      4 - FUNCTIONAL =      .2841 ESTIMATION = .28406      I
I SECOND BLANCH - VARIABLE 19 =      .8326 FORCED TOWARDS UPPER ACTIVITIES      I
-----
A 33 0 19 8 1.18997 .28406 .83260
-----

I REACH MODE      5 - FUNCTIONAL =      .1334 ESTIMATION = .13340      I
I - - - - - INTEGER SOLUTION      I
-----

```

..MPSX-V147.. EXECUTOR. MPSX RELEASE 1 PCD LEVEL 7

SOLUTION

TIME = 0.00 MINS. ITERATION NUMBER = 33

....NAME....ACTIVITY....	DEFINED AS
FUNCTIONAL	.13340	CEJ
RESTRAINTS		RPS
BOUNDS....		DECOND

- 2043

...	ROW...	AT	...	ACTIVITY...	SLACK	ACTIVITY	...	LOWER	LIMIT.	...	UPPER	LIMIT.	...	DUAL	ACTIVITY
	OBJ	BS		.13340		.13340-		NONE			NONE			1.00000	
	ROW001	BS		.14380-		.01870-		.16250-			NONE			.	
	ROW002	EQ		1.00000		.		1.00000			1.00000			.33330-	
	ROW003	EQ		1.00000		.		1.00000			1.00000			1.80000-	
	ROW004	EQ		1.00000		.		1.00000			1.00000			.40000-	
	ROW005	EO		1.00000		.		1.00000			1.00000			.96670-	
	ROW006	EQ	13330	
	ROW007	BS		3.71800		.28200		NONE			4.00000			.	
	ROW008	BS		4.56660-		.56660		NONE			4.00000-			.	
	ROW009	BS		9.53340		1.46660		NONE			11.00000			.	
	ROW010	BL		2.33330-		.		NONE			2.33330-			1.00000-	
	ROW011	BS		7.53340		3.46660		NONE			11.00000			.	

..MPSX-V1M7.. EXECUTOR. MPSX RELEASE 1 MCD LEVEL 7

SECTION 2 - COLUMNS

NUMBER	COLUMN	AT	ACTIVITY...	INPUT COST..	LOWER LIMIT.	UPPER LIMIT.
13	LAMBDA	B5	.13340	1.00000	.	1.00000
14	X001	IV	.	.	.	1.00000
15	X002	IV	1.00000	.	.	1.00000
16	X003	IV	.	.	.	1.00000
17	X004	IV	.	.	.	1.00000
18	X005	IV	.	.	.	1.00000
19	X006	IV	1.00000	.	.	1.00000
20	X007	IV	.	.	.	1.00000
21	X008	IV	1.00000	.	.	1.00000
22	X009	IV	.	.	.	1.00000
23	X010	IV	.	.	.	1.00000

..MPSX-V1M7.. EXECUTOR. MPSX RELEASE 1 MCD LEVEL 7

MIXFLOW - TIME = 0.00

 I BEST POSSIBLE SOLUTION IS NOW .13340
 I BESTORE MODE 5 - FUNCTIONAL = .13340 ESTIMATION = .13340
 I OPTIMAL INTEGER SOLUTION

APPENDIX C

Fuzzy Linear Integer Programming
Via an Exchange Heuristic, Variable
Definition Guide, User's Guide
and Sample Output


```

HEADLY M      .. HASP-II*****START JOB 1305.....NP129486.....HEADLY M
HEADLY M      .. HASP-II*****START JOB 1305.....NP129486.....HEADLY M
HEADLY M      .. HASP-II*****START JOB 1305.....NP129486.....HEADLY M

```

TIME: 15:12:48 DATE: 03/17/80 H A S P S Y S T E M L O G

```

**NP129486 JOB 04771, T=0025, R=02500, S=280, 'HEADLY M' AA
***FULLSKIPS
// EXEC SETTAF, TRAIN=TN, FORMS=16
***INCLUDE MGH01.$ONE
IEF142I - STEP WAS EXECUTED - COND CODE 0000
IEF373I STEP /SETTAF / START 80077.1512
IEF374I STEP /SETTAF / STOP 80077.1512 CPU 0MIN 00.04SEC MAIN 8K LCS OK
// EXEC FUCG
//SYSIN DD *
//DATA.FT80F001 DD UNIT=SYSDA, DSN=88EXDATA, SPACE=(3200, (50,5)),
// DISP=(NEW, PASS), DCB=(RECFM=FE, LRFCL=80, BLKSIZE=3200)
//DATA.INPUT DD *
IEF142I - STEP WAS EXECUTED - COND CODE 0000
IEF373I STEP /DATA / START 80077.1512
IEF374I STEP /DATA / STOP 30077.1512 CPU 0MIN 00.36SEC MAIN 280K LCS OK
// EXEC FUCG
//SYSIN DD *
***INCLUDE MGH01.$TWO
***INCLUDE MGH01.$THREE
//DATA.FT80F001 DD UNIT=SYSDA, DSN=88EXDATA, DISF=(OLD,DELETE)
IEF142I - STEP WAS EXECUTED - COND CODE 0000
IEF373I STEP /DATA / START 80077.1512
IEF374I STEP /DATA / STOP 80077.1512 CPU 0MIN 01.34SEC MAIN 280K LCS OK
IEF375I JOB /NP129486/ START 80077.1512
IEF376I JOB /NP129486/ STOP 80077.1512 CPU 0MIN 01.74SEC

```


[illegible]


```

** ** ** ** ** ** ** ** X6(I,J): COEFFICIENT OF THE JTH VARIABLE IN THE
** ** ** ** ITH TYPE3 FUZZY OBJECTIVE FUNCTION.
** ** ** ** ASF3(I): ASPIRATION LEVEL ITH TYPE3 FUZZY OBJECTIVE
** ** ** ** STINT(I): LOWEST ADMISSABLE VALUE OF THE ITH
** ** ** ** TYPE3 FUZZY CEJECTIF.
** ** ** ** UTINT(J): GREATEST ADMISSABLE VALUE OF THE ITH
** ** ** ** TYPE3 FUZZY CEJECTIVE.
** ** ** ** INTERNAL VARIABLES
** ** *****
** ** KOK(I,J): COEFFICIENT OF THE JTH VARIABLE IN THE ITH
** ** CONSTRAINT/OBJECTIVE FORMULATED FOR INPUT INTO
** ** THE EXCHANGE HEURISTIC.
** ** RNS(I): RIGHT-HAND SIDE VALUE OF THE ITH CONSTRAINT/
** ** OBJECTIVE FORMULATED FOR INPUT INTO THE
** ** EXCHANGE HEURISTIC.
** ** NPES: TOTAL NUMBER OF CONSTRAINTS/OBJECTIVES TO
** ** BE UTILIZED IN THE EXCHANGE HEURISTIC.
** ** NFUZZY: TOTAL NUMBER OF FUZZY FUNCTIONS.
** ** OLJCOF(J): SUM OF THE COEFFICIENTS OF THE JTH VARIABLE
** ** IN THE FUZZY FUNCTION.
** ** TECHCO(I,J): COEFFICIENT OF THE JTH VARIABLE IN THE
** ** ITH CONSTRAINT/OBJECTIVE.
** ** PVLINE(J): SUM OF THE TECECC(I,J)/AMTRES(I) RATIO
** ** FOR EACH CONSTRAINT/OBJECTIVE.

```



```

** ** ** ** ** ** ** **  ISTANT: =0, FIRST TIME PROGRAM IS EXECUTING
** ** ** **  THE FIRST SEARCH.
** ** ** **  =1, PROGRAM IS REPEATING THE FIRST SEARCH.
** ** ** **  LFITBS(K): STORE ALL VARIABLES THAT ARE PART OF THE
** ** ** **  FITBACK SOLUTION.
** ** ** **  LISTDR(J): STORE THE VARIABLES RANKED IN ORDER CF
** ** ** **  DECREASING VALUE OF RATIC(J).
** ** ** **  NNNSPRG: NUMBEF OF VARIABLES FROM THE SET OF NON-
** ** ** **  SILPCTFD PROJCTFS WHICH GAVE THE LAST ADVANTAGECUS AND
** ** ** **  FFASIBLF EXCHANGE.
** ** ** **  NSEPRO: NUMBEF OF VARIABLES FROM THE SET OF SELECTED
** ** ** **  PROJECTS WHICH GAVE THE IAST ADVANTAGEOUS AND
** ** ** **  FEASIBLE EXCHANGE.
** ** ** **  NVAFTB: NUMBEF OF VARIABLES IN THE FITBACK SOLUTION.
** ** ** **  NVARNBS: NUMBEI CF VARIABLES IN THE SET CF NON-SELECTED
** ** ** **  PROJECTS.
** ** ** **  NVARSE: NUMBR OF VARIABLES IN THE SET OF SELECTED
** ** ** **  PROJCFCTS.
** ** ** **  OPTJUN: THE VALUF OF THE MEMBERSHIP FUNCTION.
** ** ** **  SLACK(I): STORF THF SLACK VALUF FCF CONSTRAINT I.
** ** ** **  SLCKFT(I): STCRF THF SLACK VALUF FCF CONSTRAINT I
** ** ** **  DURING THF SEAPCH FOR THF FITBACK SOLUTION.
** ** ** **  *****

```



```

**      **      **      **      **      **      **      **      **      **
**      **      **      **      **      **      **      **      **      **
**      **      **      **      **      **      **      **      **      **
**      **      **      **      **      **      **      **      **      **
**      **      **      **      **      **      **      **      **      **
**      **      **      **      **      **      **      **      **      **
**      **      **      **      **      **      **      **      **      **
**      **      **      **      **      **      **      **      **      **
**      **      **      **      **      **      **      **      **      **
(RHS3(I),I=1,KTYPE3)
CARD SET # 5                      9F8.4
NUMBER OF CARDS IN SET : ITYPE1+1
FUZZY CONSTRAINT TYPE # 1
((X4(J,I),J=1,NVAR),J=1,ITYPE1)
(ASP1(I),TINT1(I),I=1,ITYPE1)
**      **      **      **      **      **      **      **      **      **
**      **      **      **      **      **      **      **      **      **
**      **      **      **      **      **      **      **      **      **
**      **      **      **      **      **      **      **      **      **
**      **      **      **      **      **      **      **      **      **
**      **      **      **      **      **      **      **      **      **
**      **      **      **      **      **      **      **      **      **
**      **      **      **      **      **      **      **      **      **
**      **      **      **      **      **      **      **      **      **
CARD SET # 6                      9F8.4
NUMBER OF CARDS IN SET : ITYPE2+1
FUZZY CONSTRAINT TYPE # 2
((X5(J,I),J=1,NVAR),J=1,ITYPE2)
(ASP2(I),TINT2(I),I=1,ITYPE2)
**      **      **      **      **      **      **      **      **      **
**      **      **      **      **      **      **      **      **      **
**      **      **      **      **      **      **      **      **      **
**      **      **      **      **      **      **      **      **      **
**      **      **      **      **      **      **      **      **      **
**      **      **      **      **      **      **      **      **      **
**      **      **      **      **      **      **      **      **      **
**      **      **      **      **      **      **      **      **      **
**      **      **      **      **      **      **      **      **      **
CARD SET # 7                      9F8.4
NUMBER OF CARDS IN SET : ITYPE3+1
FUZZY CONSTRAINT TYPE # 3
((X6(J,I),J=1,NVAR),J=1,ITYPE3)
(ASP3(I),STINT(I),I=1,ITYPE3)
**      **      **      **      **      **      **      **      **      **
**      **      **      **      **      **      **      **      **      **
**      **      **      **      **      **      **      **      **      **

```



```

READ 2, (RHS1(I), I=1, KTYPE1)
DO 100 I=1, KTYPE1
  N=M+1
  DO 200 J=1, NVAR
    ROW (M, J) = X1(I, J)
    RHS (M) = RHS1(I)
  CONTINUE
100 CONTINUE
30 H=KTYPE1
C ***
C ***
C ***
C ***
C ***

RIGID CONSTRAINT TYPE 2 (>)

IF (KTYPE2.EQ.0) GO TO 31
DO 51 IA=1, KTYPE2
  READ 2, (X2(IA, I), I=1, NVAR)
  READ 2, (RHS2(I), I=1, KTYPE2)
  DO 151 IA=1, KTYPE2
    IF (RHS2(IA).GE.0.0) GC TC 151
    RHS2(IA) = -RHS2(IA)
    DO 152 I=1, NVAR
      X2(IA, I) = -X2(IA, I)
    CONTINUE
152 CONTINUE
151 CONTINUE
DO 300 I=1, KTYPE2
  N=M+1
  DO 400 J=1, NVAR
    ROW (M, J) = X2(I, J)
    RHS (M) = RHS2(I)
  CONTINUE
300 CONTINUE
11 M=KTYPE1+KTYPE2
C ***
C ***

```

```

C *** 1, ID CONSTRAINT TYPE 3 (=) .
C ***
C ***
52 IF (KTYPE3.EQ.0) GO TO 32
DO 52 JB=1, KTYPE3
  READ 2, (X3(IB,I), I=1, NVAR)
  READ 2, (RHS3(I), I=1, KTYPE3)
DO 500 I=1, KTYPE3
  M=M+1
DO 600 J=1, NVAR
  ROW(M,J)=X3(I,J)
  RHS(M)=RHS3(I)
  M=M+1
  ROW(M,J)=-X3(I,J)
  RHS(M)=-RHS3(I)
  IF (J.EQ.NVAR) GO TO 600
  M=M-1
600 CONTINUE
500 CONTINUE
32 M=KTYPE1+KTYPE2+KTYPE3*2
C ***
C ***
C *** FUZZY CONSTRAINT TYPE 1 (<) .
C ***
C ***
IF (ITYPE1.EQ.0) GO TO 33
DO 60 ID=1, ITYPE1
  READ 2, (X4(ID,I), I=1, NVAR)
  READ 2, (ASPI(I), TINT1(I), I=1, ITYPE1)
DO 700 I=1, ITYPE1
  M=M+1
DO 800 J=1, NVAR
  ROW(M,J)=-X4(I,J)/TINT1(I)
  RHS(M)=(TINT1(I)-LSP1(I))/TINT1(I)
800 CONTINUE
300 CONTINUE

```

```

104 WRITE(60,104) (ROW(M,MN),MN=1,NVAR),RHS(M)
700 FORMAT(8F10.4)
CONTINUE
33 M=KTYPE1+KTYPE2+ KTYPE3*2+ITYPE1
C ***
C ***
C *** FUZZY CONSTRAINT TYPE 2 (>).
C ***
C ***
IF(ITYPE2.EQ.0) GO TO 34
DO 53 IF=1,ITYPE2
53 READ 2,(X5(IF,I),I=1,NVAR)
READ 2,(ASP2(I),TINT2(I),I=1,ITYPE2)
DO 900 I=1,ITYPE2
H=M+1
DO 1000 J=1,NVAR
ROW(M,J)=X5(I,J)/TINT2(I)
RHS(M)=(TINT2(I)+ASP2(I))/TINT2(I)
1000 CONTINUE
C
WRITE(6,104) (ROW(M,MN),MN=1,NVAR),RHS(M)
701 FORMAT(80,104) (RCH(M,MN),MN=1,NVAR),RHS(M)
CONTINUE
900 M=KTYPE1+ KTYPE2+ KTYPE3*2+ ITYPE1+ ITYPE2
34
C ***
C ***
C *** FUZZY CONSTRAINT TYPE 3 (=).
C ***
C ***
IF(ITYPE3.EQ.0) GO TO 35
DO 54 IG=1,ITYPE3
54 READ 2,(X6(IG,I),I=1,NVAR)
READ 2,(ASP3(I),STINT(I),UTINT(I),I=1,ITYPE3)
DO 1100 I=1,ITYPE3
M=M+1
DO 1200 J=1,NVAR
ROW(M,J)=-X6(I,J)/STINT(J)
RHS(M)=- (ASP3(I)-STINT(I))/STINT(I)
1100 CONTINUE
M=M+1

```

```

COUNT=M
ROW (M,J)=X6 (1,J)/UTINT (I)
RHS (M) = (UTINT (I) +ASP3 (I))/UTINT (I)
IF (J.EQ.NVAR) GO TO 1200
M=M+1
1200 CONTINUE
KKJ=COUNT-1
C WRITE (6,104) (ROW (KKJ,MK),MK=1,NVAR),RHS (KKJ)
WRITE (80,104) (ROW (KKJ,MK),MK=1,NVAR),RHS (KKJ)
KKJ=KKJ+1
C WRITE (6,104) (ROW (KKJ,MK),MK=1,NVAR),RHS (KKJ)
WRITE (80,104) (ROW (KKJ,MK),MK=1,NVAR),RHS (KKJ)
KKJ=0
1100 CONTINUE
1 FORMAT (7I5)
2 FORMAT (9F8.4)
35 KJK=KTYPE1+KTYPE2+KTYPE3*2
DC 1124 LMN=1,KJK
C WRITE (6,104) (ROW (LMN,J),J=1,NVAR),RHS (LMN)
1124 WRITE (80,104) (ROW (LMN,J),J=1,NVAR),RHS (LMN)
STOP
END

```

TFD= 746

OBJECT CODE= 12624 BYTES,ARRAY AREA= 18600 BYTES,TOTAL AREA AVAILABLE= 17817

NUMBER OF ERRORS= 0, NUMBER OF WARNINGS= 0, NUMBER OF EXTENSIONS=

0.21 SEC,EXECUTION TIME= 0.06 SEC,

```

COMMON/COM1/ISELTD( 50),NSELTD( 50),ISELEC( 3),NSELEC( 3),NPRCNS,NP
+POSE,IOUTPT
COMMON/COM2/GNOBFU(2),CBJCO?(50)
COMMON/COM3/AMTAVL(30),AMTNEO(30),SLACK(30),TFCHCO(30,50),NRES,IFE
IASB
COMMON/COM4/IMPDS(6),NRSFEC,NSEFPO
COMMON/COM5/LISTPR( 50),NVARNS,NVARSE
COMMON/COM21/INDLX,FACEFU,NCOUNT
COMMON/COM25/X
COMMON/COM24/IXX
COMMON/COM22/FRES
COMMON/COM23/GJJFUN
COMMON/COM38/FCONST
COMMON/COM39/NFUZZY
COMMON/COM40/IMARK
COMMON/COM41/KFLAG,TRIAL,JFLAG
DIMENSION NS2ARY(50),X(50),PRES(50,50),FCONST(50)
DIMENSION AMTRES(30),IFITES(50),KATIC(50),RVALUE(50),SLCKFT(30)
DIMENSION IA(50)
NVAESE=0
NVAENS=0
WRITE(6,840)

C ***
C *** READ DATA
READ(80,850) NVAR,NRES,NFUZZY
IOUTPT=0
IF(NVAR.EQ.0) STOP
DO 2 II=1,NRES
READ(80,860) (TECHCO(II,JJ),JJ=1,NVAR),AMTRES(II)
2 CONTINUE
DO 622 I=1,NVAR
ISELTD(I)=0
NSELTD(I)=0
622 CONTINUE
DO 1134 JJ=1,NVAR
OBJECT(JJ)=0.0

```

```

DO 1133 I=1,NFUZZY
  OBJCOF(JJ)=OBJCOF(JJ)+TECHCO(I,JJ)
1133 CONTINUE
1134 CONTINUE
DO 133 I=1,NFUZZY
  DO 134 JJ=1,NVAR
    PRES(I,JJ)=TECHCO(I,JJ)
    FCONST(I)=AMTRES(I)
134 CONTINUE
133 CONTINUE
  ISFART=0
  C *** CALCULATE RVALUE FOR EACH VARIABLE
  DO 18 J=1,NVAR
    RVALUE(J)=0.0
    DO 19 I=1,NRES
      IF (AMTRES(I).EQ.0.0) GC TC 18
      RVALUE(J)=RVALUE(J)+TECHCO(I,J)/AMTRES(I)
18 CONTINUE
  C ***
  C *** CALCULATE RATIO
  RFLAG=0
  DO 20 J=1,NVAR
    LISTDR(J)=J
    IF (RVALUE(J).EQ.0.0) GC TC 1120
    RATIO(J)=OBJCOF(J)/RVALUE(J)
    GC TO 20
1120 RATIO(J)=0.0
  20 CONTINUE
  C ***
  C *** BANK PROJECTS BASED ON VALUE OF RATIO(J) - HIGHEST=1ST, ...
  DO 40 ID=2,NVAR
    IID=10
  30 IF (RATIO(IID).LE.RATIO(IID-1)) GC TC 40
  C ***
  C *** INTERCHANGE RATIO
  RATIO(IID)=RATIO(IID-1)

```



```

      RATIO(IID-1)=RATIO(IID)
      PATIO(IID)=TRATIO

C *** INTERCHANGE PROJECT NUMBER
C *** LIST=LISTDR(IID-1)
      LISTDR(IID-1)=LISTDR(IID)
      LISTDR(IID)=LIST
      IID=IID-1
      IF (IID-1) 40,40,30
40 CONTINUE

C *** DETERMINE THE INITIAL SOLUTION WITHOUT FITBACK.
C *** 181 NT=NEUZZY+1
      NCOUNT=0
      SUM=0.0
      NVARNS=0
      DO 111 KJ=1,NVAR
111   IA(KJ)=0
      NTIM=(2**NVAR)-1
      DO 201 MP=1,NTIM
      INDEX=0
      NVARSE=0
      INDEX=0
      JJ=1
35   JT=LISTDR(JJ)
      IA(JT)=IA(JT)+1
      IF (IA(JT).NE.2) GO TO 50
      IA(JT)=0
      JJ=JJ+1
      GO TO 35
50   DO 778 KLJ=1,NVAR
      J=LISTDR(KLJ)
      IF (IA(J).EQ.0) GO TO 165
      NVARSE=NVARSE+1
      LISTDR(NVARSE)=LISTDR(KLJ)

```

```

GO TO 778
165 INDEX=INDEX+1
NSELTD(INDEX)=LISTDR(KIJ)
778 CONTINUE
DO 80 I=NT,NRES
SUM=0.0
DO 60 IND=1,NVAR
J=LISTDR(IND)
IF(IA(J).EQ.0) GO TO 60
SUM=SUM+TECHCO(I,J)
GO TO 60
60 CONTINUE
IF(SUM.LE.AMRES(1)) GO TO 70
NCOUNT=0
GO TO 201
70 NSLACK=0
DO 110 IJK=NT,NRES
AMTSUD=0.0
IF(NVARSE.EQ.0) GO TO 201
DO 1100 IM=1,NVARSE
KND=ISELTD(IM)
1100 AMTSUD=AMTSUD+TECHCO(IJK,KND)
SLACK(IJK)=AMRES(IJK)-AMTSUD
IF(SLACK(IJK).GE.0.0) GC TC 1121
NCOUNT=0
GO TO 201
1121 NSLACK=NSLACK+1
1110 CONTINUE
IF(NSLACK.EQ.(NRES-NFUZZY)) GO TO 85
GO TO 80
85 DO 1190 JT=1,NFUZZY
AMTSUD=0.0
DO 1199 IQ=1,NVARSE
IT=ISELTD(IQ)
1199 AMTSUD=AMTSUD+TECHCO(JT,IT)
SLACK(JT)=AMRES(JT)-AMTSUD

```

```

IF(SLACK(JT).GE.0.0) GO TO 1198
RCOUNT=0
GO TO 201
1198 CONTINUE
80 CONTINUE
IF(HVARSE.EQ.0) GO TO 201
GO TO 1185
201 CONTINUE
C *** FILL THE VECTORS OF SELECTED AND NON-SELECTED VARIABLES
1185 WRITE(6,910) (ISELTD(I),II=1,NVARSF)
NVARNS=INDEX
WRITE(6,924) (NSELTD(IJ),IJ=1,NVAFNS)
924 FORMAT(IX,'THE SET OF NON-SELECTED VARIABLES IS ',16I5)
C ***
C *** CALCULATE THE MEMBERSHIP FUNCTION AS AN INITIAL SOLUTION
OBJFUN=0.0
CALL OBJ
IF(IMARK.EQ.10) GO TO 95
WRITE(6,920) OBJFUN
C *** CALCULATE SLACK FOR EACH CONSTRAINT
95 DO 110 I=1,NRES
AMTSUD=0.
DO 100 IS=1,NVARSE
IND=ISELTD(IS)
100 AMTSUD=AMTSUD+TECHCC(I,IND)
SLACK(I)=AMTRES(I)-AMTSUD
IF (LOUPI.EQ.0) GO TO 110
IF(IMARK.EQ.10) GO TO 110
WRITE(6,925) I,SLACK(I)
110 SLCKFT(I)=SLACK(I)
C ***
C *** RANK PROJECTS IN BOTH SETS OF SELECTED AND NON - SELECTED
C *** PROJECTS BASED ON THEIR OBJECTIVES' COEFFICIENTS VALUE
195 CALL RANK

```

```

C ***
C *** IF (ISTART.EQ.1) GC TO 150
C *** FIND FITBACK SOLUTION
DO 220 IP=1,NVARSE
220 LFTTRS(IP)=ISELTD(IP)
NVAFTB=NVARSE
FTOBFU=OPJFUN
NCOUNT=0
DO 140 J=1,NVARNS
IFLAG=0
INDEX=NSELTD(J)
DO 130 II=1,NRES
IF(SLCKFT(II).LT.TECHCC(II,INDEX)) GC TO 130
IFLAG=IFLAG+1
130 CONTINUE
IF(IFLAG.NE.NRES) GO TO 140
C *** FOUND A VARIABLE FOR FITBACK SOLUTION
C *** UPDATE SLACK VALUE AND CALCULATE
C *** THE MEMBERSHIP FUNCTION
DO 135 IK=1,NRES
SLCKFT(IK)=SLCKFT(IK)-TECHCC(IK,INDEX)
NVAFTL=NVAFTB+1
IF(ILLS(NVAFTB)=NSELTD(J))
CALL FTOBJ
NCOUNT=NCOUNT+1
140 CONTINUE
IF(NVAFTB.EQ.NVARST) GC TO 141
WRITE(6,930) (LFTTRS(II),II=1,NVAFTB)
WRITE(6,920) FTOBFU
WRITE(6,932)
GO TO 150
141 WRITE(6,142)
142 FORMAT(5X,'NO FITBACK SOLUTION CAN BE FOUND TO IMPROVE THE INITIAL
1 SOLUTION')
FTOBFU=-9999.
C ***

```

```

C *** 2/1 EXCHANGE
150 NPROBS=2
    NPROSE=1
    NPROFT=0
    GNOBFU(1)=0.
    NS1LST=NVARNS-1
    IF (NS1LST.EQ.0) GC TC 212
    DO 152 INS=1, NS1LST
    NS2ARY(INS)=NVARNS
    DO 210 JS1=1, NVARSE
    ISELEC(1)=ISELTD(JS1)
    DO 190 NS1=1, NS1LST
    NS1NP=NS1
    NSELEC(1)=NSELTD(NS1)
    NS2LST=NS2ARY(NS1)
    NS2ST=NS1+1
    IF (NS2ST.GT.NS2LST) GC TC 210
    IFLAG=0
    DO 170 NS2=NS2ST, NS2LST
    NS2NP=NS2
    NSELEC(2)=NSELTD(NS2)
    C *** CHECK MEMBERSHIP FUNCTION
    CALL ACHVMT
    IF (GNOBFU(2).GE.GNOBFU(1)) GO TO 160
C *** 2/1 EXCHANGE NOT ADVANTAGEOUS
C *** STORE PROJECT NUMBERS SO THAT THEY WILL NOT BE USED IN THE 1/1 EXC
    IF (NS1.NE.1.CH.NEGFT.NE.0) GO TO 155
    NPROFT=JS1
155 IFLAG=1
    IF ((NS2-NS1)-2) 200,200,180
    C *** CHECK FEASIBILITY
160 CALL FCASEL
    IF (LFEASE.NE.NRFS) GO TO 170
    CALL IMPVNT
    IFLAG=1
    IF ((NS2-NS1)-2) 200,200,180

```

```

170 CONTINUE
180 NS2LST=NS2NP-1*IFLAG
    IEND=NVARNS-1
    DO 185 I=NS1,IEND
    IF(NS2ARY(I).NE.NS2LST) GO TO 185
    NS2ARY(I)=NS2LST
185 CONTINUE
190 CONTINUE
200 IDIF=(NS2NP-NS1NP)*IFLAG
    NS1LST=NS1NP-((2-IDIF)*IFLAG)
    IF(NS1LST.EQ.0) GO TO 212
    IF(NS1NP.NE.NS1LST) GO TO 210
    NS2ARY(NS1NP)=NS2NP+((1-IDIF)*IFLAG)
210 CONTINUE
C ***
C *** 1/1 EXCHANGE
212 NPRONS=1
    JS1LST=NVARCE+NPROFT
    IF(JS1LST.GT.NVARSP)JS1LST=NPROFT-1
    DO 240 NS1=1,NVARNS
    NSELEC(1)=NSELTD(NS1)
    IFLAG=0
    DO 221 JS1=1,JS1LST
    ISELFC(1)=ISELTD(JS1)
    JS1NP=JS1
    CALL ACHVMT
    IF(GHODFU(2).GE.GNCBFU(1)) GO TO 215
    IFLAG=1
    GO TO 230
215 CALL FEASEL
    IF(IFAASH.NE.NRES) GC TC 221
    CALL IMPVNT
    IFLAG=1
    GO TO 230
221 CONTINUE
C ***

```

```

C *** 1/2 EXCHANGE
230 JS1LST=JSINP-1*YFLAG
    IF (JS1LST.EQ.0) GC 1C 250
240 CONTINUE
250 IF (GNORFU(1)-LE.0.) GO TO 290
    OBJFUN=OBJFUN+GNORFU(1)
C *** FOUND AN EXCHANGE TO IMPROVE THE MEMBERSHIP FUNCTION
    IF (IOUTPT.EQ.0) GC TO 260
    WRITE(6,935)
    CALL OBJ
    WRITE(6,920) OBJFUN
    WRITE(6,940)

C *** MAKE THE EXCHANGE
260 CALL EXCHGE
C ***
C *** REPEAT FIRST SEARCH
    ISTART=1
    GO TO 95
290 IF (ISTART.EQ.0) GC 1C 300
    WRITE(6,935)
    WRITE(6,950) (ISFITD(II),II=1,NVARS)
    CALL OBJ
    WRITE(6,920) OBJFUN
    GO TO 305
300 WRITE(6,955)
305 IF (NVARNS.LT.3) GO TO 480
C *** SECOND SEARCH (3/1,3/2,3/3 EXCHANGES)
    WRITE(6,960)
    NPRONS=3
    GNORFU(1)=0.
    JS1LST=NVARNS-2
    ISTEP2=0
    ISTEP3=0
    DO 430 NS 1=1,NS1LST
C ***

```

```

C *** 3/1 EXCHANGE
DO 310 IIP=1,3
310 NSELEC(IIP)=NSELTD(NS1+IIP-1)
NPROSF=1
IFLAG=1
DO 330 JS1=1,NVARSE
JS1NP=JS1
JSELEC(1)=ISELTD(JS1)
C ***
C *** CHECK FOR PROFITABILITY
CALL ACHVMT
IF(GNORFU(2).GF.GNORFU(1)) GO TO 320
IFLAG=0
GO TO 340
C ***
C *** CHECK FOR FEASIBILITY
320 CALL FEASBL
IF(IFEASE.NE.NRES) GO TO 330
C ***
C **** STORE IMPROVED SOLUTION TEMPORARILY
CALL IMPVNT
330 CONTINUE
340 IF(JS1NP.LP.(2-IFLAG).CH.ISTP32.EQ.1) GO TO 430
C ***
C *** 3/2 EXCHANGE
JS1LST=JS1NP-2+IFLAG
JS2LST=JS1NP-1+IFLAG
JSLIMT=JS1NP
NPROSE=2
DO 370 JS1=1,JS1LST
ISELEC(1)=ISELTD(JS1)
JS2ST=JS1+1
IF(JS2ST.GT.JS2LST) GO TO 380
IFLAG=0
DO 360 JS2=JS2ST,JS2LST
JS2HP=JS2

```



```

C ***
C ***
ISELEC(2)=ISELTD(JS2)
CHECK FOR PROFITABILITY
CALL ACHVMT
IF(GNOBFU(2).GE.GNOBFU(1)) GO TO 350
IF(JS1.NE.1.CR(JS2.NE.2) GO TO 345
ISTPJ2=1
GO TO 430
345 IF((JS2-JS1).LE.2) GO TO 380
IFLAG=1
GO TO 365
C ***
C ***
CHECK FOR FEASIBILITY
350 CALL FEASBL
IF(IFLASE.NE.NRES) GO TO 360
C ***
C ***
STORE IMPROVED SOLUTION TEMPORARLY
CALL IMPVNT
CONTINUE
360 JS2LST=JS2NP-1*IFLAG
370 CONTINUE
380 IF(NVAPSE.LT.3.CR(JS1LST.EC.1-OR.ISTP33.E2.1) GO TO 430
C ***
C ***
3/3 EXCHANGE
NPROSF=1
JS1LST=JS1LST-1
JS2LST=JS2LST-1
JS3LST=JS3LST-1
DO 420 JS1=1,JS1LST
ISELEC(1)=ISELTD(JS1)
JS2ST=JS1+1
DO 410 JS2=JS2ST,JS2LST
ISELEC(2)=ISELTD(JS2)
JS3ST=JS2+1
DO 400 JS3=JS3ST,JS3LST
ISELEC(3)=ISELTD(JS3)
C ***

```

```

C *** CHECK FOR PROFITABILITY
CALL ACHVMT
IF(GHOFU(2).GE.GHOFU(1)) GO TO 390
IF(JS1.NE.1.CR.JS2.NE.2.OR.JS3.NE.3) GO TO 385
ISTP33=1
GO TO 430
385 IF((JS3-JS2)-2) 430,430,410
C ***
C *** CHECK FOR FEASIBILITY
390 CALL FEASEL
IF(IFEASD.NE.NRES) GO TO 400
C ***
C *** STORE IMPROVED SOLUTION TEMPORARILY
CALL IMPVNT
400 CONTINUE
410 CONTINUE
420 CONTINUE
430 CONTINUE
IF(GHOFU(1).GT.0.) GO TO 470
WRITE(6,970)
GO TO 480
470 WRITE(6,980)
CALL OBJ
480 IF(OBJFUN.GE.FTOBFU) GO TO 490
WRITE(6,990) (LFITB(I),II=1,NVAFTE)
WRITE(6,920) FTOBFU
GO TO 500
490 WRITE(6,990) (ISELTD(II),II=1,NVASE)
WRITE(6,920) OBJFUN
500 STOP
840 FORMAT(1E0,36X,59H*** HEURISTIC ALGORITHM FOR LARGE 0 - 1 LINEAR P
+LOGRAMS ***)
850 FORMAT(16I5)
860 FORMAT(8F10.4)
910 FORMAT(1H0,64HTHE INITIAL SOLUTION (WITHOUT FITBACK) IS COMPOSED O
+ F PRODUCTS : ,50I4///)

```

```

920  FORMAT(1H0,41HTHE VALUE OF THE MEMBERSHIP FUNCTION IS :,F8.2///)
930  FORMAT(1H0,46HTHE FITBACK SOLUTION IS COMPOSED OF PROJECTS :,50I4)
932  FORMAT(1H0,45X,41H*** FIRST SEARCH (2/1 - 1/1 EXCHANGE) ***//)
935  FORMAT(1H0,73HTHE FIRST SEARCH HAS FOUND AN EXCHANGE TO IMPROVE TH
      +E MEMBERSHIP FUNCTION)
940  FORMAT(1H0,46X,40HREPEAT FIRST SEARCH (2/1 - 1/1 EXCHANGE)///)
950  FORMAT(1H0,38HTHE SOLUTION IS COMPOSED OF PROJECTS :,50I4)
955  FORMAT(1H0,73HTHE FIRST SEARCH HAS FOUND NO EXCHANGE TO IMPROVE TH
      +E MEMBERSHIP FUNCTION)
990  FORMAT(1H1,44HTHE FINAL SOLUTION IS COMPOSED OF PROJECTS :,50I4)
925  FORMAT(1H0,26HTHE SLACK FOR CONSTRAINT :,I4,5H IS :,F8.2)
960  FORMAT(1H0,42X,1  *** SECOND SEARCH (3/1 3/2 3/3 EXCHANGE) ***
      +)
970  FORMAT(1H0,76HTHE SECOND SEARCH HAS FOUND NO EXCHANGES TO IMPROVE
      +THE MEMBERSHIP FUNCTION.)
930  FORMAT(1H0,75HTHE SECOND SEARCH HAS FOUND AN EXCHANGE TO IMPROVE T
      +HE MEMBERSHIP FUNCTION.)
      FND

```

```

SUBROUTINE RANK
  DIMENSION X(50)
  COMMON/COM1/ISELTD(50),NSELTD(50),ISELEC(3),NSELEC(3),NPRONS,NPR
  +OSE,JOBTPT
  COMMON/COM2/GMOBEU(2),OBJCOF(50)
  COMMON/COM5/IISTDR( 50),NVARNS,NVAFSE
  COMMON/COM25/X

```

```

C ***
C *** RANK PROJECTS IN SET OF NOV - SELECTED PROJECTS (NSELTD) BASED
C *** ON OBJCO VALUE
      IF(NVAFNS.FO. 1) GO TO 21
      DO 20 ID=2,NVARNS
      IID=ID
10  IED1=ISELTD(IID-1)

```

```

IND2=NSELTD(IID)
IF (OBJCOF(IND2).LE.OBJCOF(IND1)) GO TO 20
NSELTT=NSELTD(IID-1)
NSELTD(IID-1)=NSFLTD(IID)
NSFLTD(IID)=NSELTT
IID=IID-1
IF (IID-1) 20,20,10
20 CONTINUE
21 CONTINUE
C ***
C *** PANK PROJECTS IN SET OF SELECTED PROJECTS (ISELTD) BASED ON
C *** OBJCO VALUE
IF (NVARSE.EQ.1) GO TO 41
DO 40 I2=2,NVARSE
IID=ID
30 IND1=ISELTD(IID-1)
IND2=ISELTD(IID)
IF (OBJCOF(IND2).GE.OBJCOF(IND1)) GO TO 40
ISELTT=ISELTD(IID-1)
ISELTD(IID-1)=ISELTD(IID)
ISELTD(IID)=ISELTT
IID=IID-1
IF (IID-1) 40,40,30
40 CONTINUE
41 CONTINUE
IF (IOUTPT.EQ.0) GO TO 70
WRITE(6,50) (NSELTD(II),II=1,NVARNS)
WRITE(6,60) (ISELTD(II),II=1,NVARSE)
50 FORMAT(1H0,60THE SET OF NON - SELECTED PROJECTS IS COMPOSED OF PR
+OBJECTS :,50I4)
60 FORMAT(1H0,54THE SET OF SELECTED PROJECTS IS COMPOSED OF PROJECTS
+ :,50I4///)
70 RETURN
END

```

```

SUBROUTINE IMPVNT
COMMON/COM1/ISELTD(50),NSFLTD(50),ISELEC(3),NSELEC(3),NPRONS,NPR
+OSE,IOUTPT
COMMON/COM2/GNOBFU(2),CEJCOF(50)
COMMON/COM4/IMPDSO(6),NNSERO,NSFFERO
GNOBFU(1)=GNOBFU(2)
DO 10 IP=1,NPRONS
10 IMPDSO(IP)=NSELEC(IP)
DO 20 IP=1,NEROSE
20 IMPDSO(NPRONS+IP)=ISELEC(IP)
NNSPRG=NPRONS
NSFFRG=NSFFERO
IF (IOUTPT-EC.0) GO TO 40
WRITE(6,30)
30 FORMAT(1H0,48H*** THE EXCHANGE IS BOTH PROFITABLE AND FEASIBLE)
40 RETURN
END

```

```

SUBROUTINE FEASBL
COMMON/COM1/ISELTD(50),NSELTD(50),ISELEC(1),NSELEC(3),NPRONS,NPR
+OSE,IOUTPT
COMMON/COM3/AMTAVL(30),AMTNEO(30),SLACK(30),TECHCC(30,50),NRES,IFF
+ASB
IFEASB=0
DO 40 I=1,NRES
  AMTNEO(I)=0.
  DO 10 IP=1,NPRONS
    INDX=NSELEC(IP)
10    AMTNEO(I)=AMTNEO(I)+TECHCC(I,INDX)
    DO 22 IP=1,NPROSE
      INDY=ISELEC(IP)
22    AMTAVL(I)=SLACK(I)+TECHCO(I,INDY)
      IF(I.EQ.NRES.AND.AMTNEO(I).GT.AMTAVL(I)) GO TO 40
      IF(IOUTPT.EQ.0) GO TO 70
70    IF(AMTNEO(I).GT.AMTAVL(I)) GO TO 40
      IFEASE=IFEASE+1
40    CONTINUE
      RETURN
      END

```

```

SUBROUTINE EXCHGE
COMMON/COM1/ISELFD(50),NSELTD(50),ISELEC(3),NSEIEC(3),NPRONS,NPR
+OPE,IOUTPT
COMMON/COM4/IMPDSO(6),NNSPRO,NSEPRO
COMMON/COM5/LISTDR(50),NVARPS,NVARESE

C ***
C *** MAKE THE EXCHANGE IN SET OF NON - SELECTED PROJECTS
ISST=1
INDEX=0
DO 30 IP=1,NVARNS
IF (ISST.GT.NNSPRC) ISST=NNSPRO
DO 10 IS=ISST,NNSPRO
IF (NSELTD(IP).EQ.IMEDSC(IS)) GO TO 20
INDEX=INDEX+1
LISTDR(INDEX)=NSELTD(IP)
GO TO 30
10 CONTINUE
20 ISST=IS+1
30 CONTINUE
IF (INDEX.EQ.0) GO TO 45
DO 40 IP=1,INDEX
40 NSELTD(IP)=LISTDR(IP)
45 DO 50 IP=1,NSEPRO
50 NSELTD(INDEX+IP)=IMPDSO(NNSPRO+IP)
NVARNS=NVARNS+NSEPRC-NASPRO

C ***
C *** MAKE THE EXCHANGE IN SET OF SELECTED PROJECTS
ISST=1
INDEX=0
DO 80 IP=1,NVARESE
IF (ISST.GT.NSEPRC) ISST=NSEPRO
DO 60 IS=ISST,NSEPRO
IF (NSELTD(IP).EQ.IMEDSC(NNSPRC+IS)) GO TO 70
INDEX=INDEX+1
LISTDR(INDEX)=NSELTD(IP)
GO TO 80

```

```

60 CONTINUE
70 ISST=IS+1
80 CONTINUE
  IF (INDIX.EQ.0) GO TO 95
  DO 90 IP=1,INDEX
    90 ISELTD(IP)=LISTDR(IP)
  95 DO 100 IP=1,NNSPRO
    100 ISELTD(INDEX+IP)=IMPD50(IP)
  NVARSE=NVARSE+NNSPRO-NSEPRO
  RETURN
  END

SUBROUTINE ACHVMT
  DIMENSION FRES(50,50),FCONST(50),FOBJ(50),TEST(50),IBASE(50,50),
  1 NBASE(50,50)
  COMMON/COM1/ISELTD(50),NSELTD(50),ISELEC(3),NSELEC(3),NPRONS,NPR
  +OSE,IOUTPT
  COMMON/COM2/GNOBFU(2),OBJCOF(50)
  COMMON/COM5/LISTDR(50),NVARNS,NVARSE
  COMMON/CON24/IXX
  COMMON/CON22/FRES
  COMMON/CON34/NCOUNT
  COMMON/CON38/FCONST
  COMMON/CON39/NFUZZY
  COMMON/CON41/KFLAG,TRIAL,JFLAG
  COMMON/CON40/IMARK
  COMMON/CON23/OBJFUN
  GNOBFU(2)=0.0
  KFLAG=1
  DO 134 I=1,NFUZZY
    FOBJ(I)=0.0
  DO 145 IP=1,NVARSE
    145 ISELTD(IP)
  150 FOBJ(I)=(-FRES(I,INDX))+FOBJ(I)

```



```

145 CONTINUE
   FOBJ(I)=FCONST(1)+FOBJ(I)
   DO 135 IP=1,NPRONS
     INDX=NSELEC(IP)
     NBASE(I,IP)=NSELEC(IP)
135   FOBJ(I)=FOBJ(I)+(-FRFS(I,INDX))
   DO 136 IIP=1,NPROSE
     INDX=ISELEC(IIP)
     IBASE(I,IIP)=ISELEC(IIP)
136   FOBJ(I)=FOBJ(I)-(-FRFS(I,INDX))
     TEST(I)=FOBJ(I)
134 CONTINUE
   DO 83 I=2,NFUZZY
     J=I-1
     IF(TEST(I).LT.0.0) GO TO 83
     IF(TEST(I).GT.1.0) TEST(I)=1.0
     IF(TEST(I).LE.TEST(J)) GO TO 183
     USOFAR=TEST(J)
   DO 1283 IP=1,NPROSE
1283   ISELEC(IP)=IBASE(J,IP)
   DO 1383 IP=1,NPRONS
1383   NSELEC(IP)=NBASE(J,IP)
   GO TO 83
183 USOFAR=TFST(I)
   DO 283 IP=1,NPROSE
283   ISELEC(IP)=IBASE(I,IP)
   DO 383 IP=1,NPRONS
383   NSELEC(IP)=NBASE(I,IP)
83 CONTINUE
34   GNOBFU(2)=USOFAR-OBJFUN
   IF(GNOBFU(2).LE.0.0) GNOBFU(2)=0.0
35   IF(IOUTPT.EQ.0) GO TO 80
   WRITE(6,30) (ISELEC(I),I=1,NPROSE)
30   FORMAT(1H0,65H THE EXCHANGE UNDER CONSIDERATION CONSISTS OF REPLAC
     +65 PROJECTS :,3I4)
   WRITE(6,40) (NSELEC(I),I=1,NPRONS)
40   FORMAT(1H+,79X,14H BY IFCJECTS :,3I4)
51   WRITE(6,50) GNOBFU(2)

```

```

50 FORMAT(9X, ' THE GAIN IN THE MEMBERSHIP FUNCTION WOULD BE :', F3.2)
80 RETURN
END

```

```

SUBROUTINE OBJ
  DIMENSION FRES(50,50), FCCNST(50), FOBJ(50), YLAMDA(50)
  COMMON/COM1/ISELTD(50), NSELTD(50), ISELEC(3), NSELEC(3), NPRONS,
  INPROSE, IOUTPT
  COMMON/CON22/FRES
  COMMON/CON23/OBJFUN
  COMMON/COM5/LISTR(50), NVARNS, NVARSE
  COMMON/CON38/FCONST
  COMMON/CON39/NFUZZY
  COMMON/CON40/IMARK
  COMMON/CON41/KFLAG, TRIAL, JFLAG
  DO 134 I=1, NFUZZY
    FOBJ(I)=0.0
  DO 135 IP=1, NVARSE
    INDY=ISELTD(IP)
135  FOBJ(I)=(-FRES(I, INDX))+FOBJ(I)
    YLAMDA(I)=FCONST(I)+FCEJ(I)
    IF (YLAMDA(I).LT.0.0) GO TO 139
    IF (YLAMDA(I).GT.1.0) YLAMDA(I)=1.0
    TEST=YLAMDA(I)
    IF (I.EQ.1) GO TO 138
    IF (TEST.LT.BEST) DFST=YLAMDA(I)
    GO TO 134
138  BEST=TEST
134  CONTINUE
    OBJFUN=BFST
    IMARK=0
    GO TO 152
139  WRITE(6,140)
140  FORMAT(5X, 'THE INITIAL SOLUTION IS INFEASIBLE')
    IMARK=10

```

```

          OBJFIR=0.0
152  RETURN
      END

SUBROUTINE FTCEJ
DIMENSION FRES(50,50),FCONST(50),OBJ(50),LOBJ(50)
COMMON/CON21/INDEX,FTCEFU,MCOUNT
COMMON/CON22/FRES
COMMON/CON34/NCOUNT
COMMON/CON2/GNOBFU(2),CBJCOF(50)
COMMON/CON38/FCONST
COMMON/CON39/NFUZZY
IF(MCOUNT.GT.0) GC TO 932
DO 934 IF=1,NFUZZY
  OBJ(IF)=FCONST(IF)+FTCEFU
DO 935 IG=1,NFUZZY
  LOBJ(IG)=OBJ(IG)+FRES(IG,INDEX)
  TEST=LOBJ(IG)
  IF(IG.EQ.1) GC TO 938
  IF(TEST.LT.BEST) BEST=LOBJ(IG)
  GO TO 935
938 BEST=TEST
935 CONTINUE
  FTCEFU=BEST
  RETURN
932 DO 1034 JF=1,NFUZZY
1034 OBJ(JF)=FTCEFU
DO 1035 JG=1,NFUZZY
  LOBJ(JG)=OBJ(JG)+FRES(JG,INDEX)
  TEST=LOBJ(JG)
  IF(JG.EQ.1) GC TO 1038
  IF(TEST.LT.BEST) BEST=LOBJ(JG)
  GO TO 1035
1038 TEST=TEST
1035 CONTINUE

```

FTOBFU=BEST
RETURN
END

D= 104717

OBJECT CODE= 31024 BYTES,ARRAY AREA= 40164 BYTES,TOTAL AREA AVAILABLE

NUMBER OF ERRORS= 0, NUMBER OF WARNINGS= 0, NUMBER OF EXT

0.42 SEC,EXECUTION TIME= 0.84 SEC,

*** HEURISTIC ALGORITHM FOR LARGE 0 - 1 LINEAR PROGRAMS ***

THE INITIAL SOLUTION (WITHOUT FITBACK) IS COMPOSED OF PROJECTS: 6 2 8
 THE SET OF NON-SELECTED VARIABLES IS 7 4 5 1 3 10 9

THE VALUE OF THE MEMBERSHIP FUNCTION IS: 0.13

NO FITBACK SOLUTION CAN BE FOUND TO IMPROVE THE INITIAL SOLUTION

THE FIRST SEARCH HAS FOUND NO EXCHANGE TO IMPROVE THE MEMBERSHIP FUNCTION

**** SECOND SEARCH (3/1 3/2 3/3 EXCHANGE) ****

THE SECOND SEARCH HAS FOUND NO EXCHANGES TO IMPROVE THE MEMBERSHIP FUNCTION

THE FINAL SOLUTION IS COMPOSED OF PROJECTS: 8 6 2

THE VALUE OF THE MEMBERSHIP FUNCTION IS: 0.13

APPENDIX D

Sample Data Input

